CS 453/698: Software and Systems Security

Module: Hardware & Mobile Security

Lecture: Android & ARM TrustZone

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Reminders & Recap

Reminders:

- A4 is released
 - Due July 25th
- Send your research project proposals to Meng and me!

Recap – last time we covered:

Intro to Trusted Execution Environments (TEE)

- Separate, isolated, and minimal execution environmet
- Enabled as a part of the CPU arch. itself (not a separate external module)

Intel SGX

- User-space TEE → enclaves
- Architecture details → Isolation, life cycle, address translation, attestation

Today

Continue: Hardware and Mobile Security

Different TEE architecture and real-world use case

ARM TrustZone

- System-level or "split world" architecture
- Overview
- Architectural details

Android OS

- How it uses TrustZone (particular focus on KeyStore)
- Other security features

ARM Processors

A few family of CPUs provided by ARM

ARM Cortex-A family

- Application processors
- Supports OS and high-performance apps
- This is the CPU in smart phones, smart tv

ARM Cortex-R family

- Real-time processors with highperformance and reliability
- Support real-time processing and mission-critical control

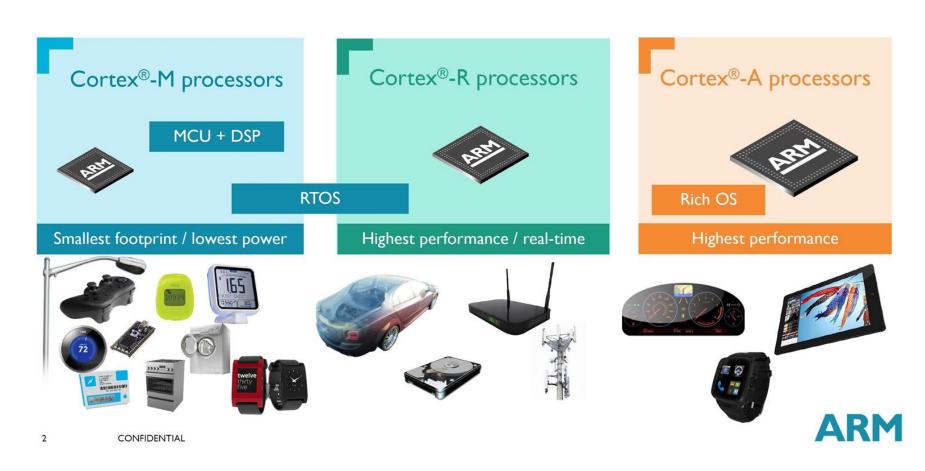
ARM Cortex-M family

- Microcontrollers
- Cost-sensitive, SoC, low-power processing

ARM Processors

A few family of CPUs provided by ARM

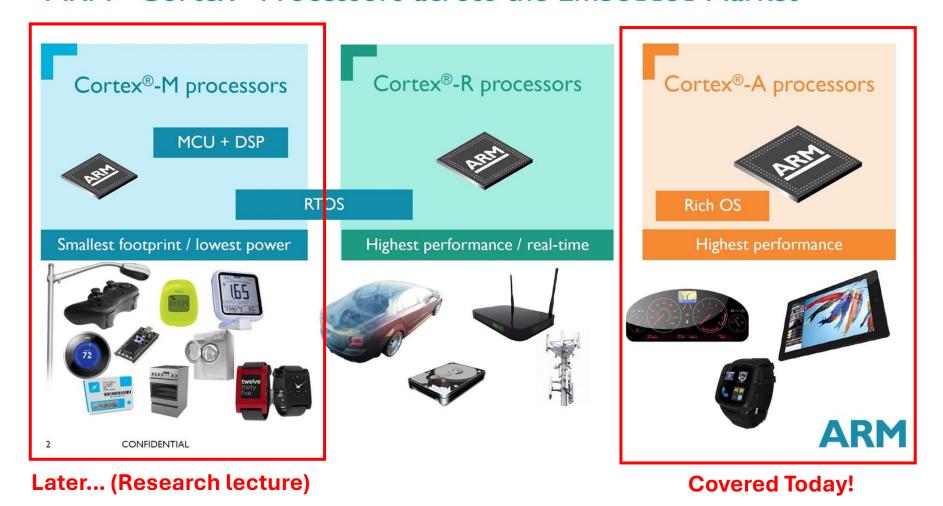
ARM® Cortex® Processors across the Embedded Market



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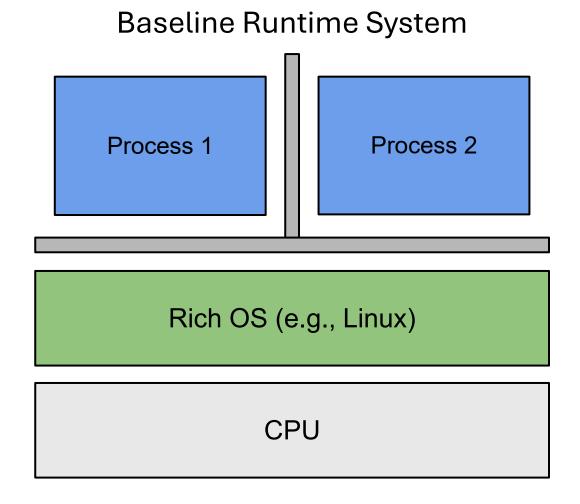
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What is TrustZone?

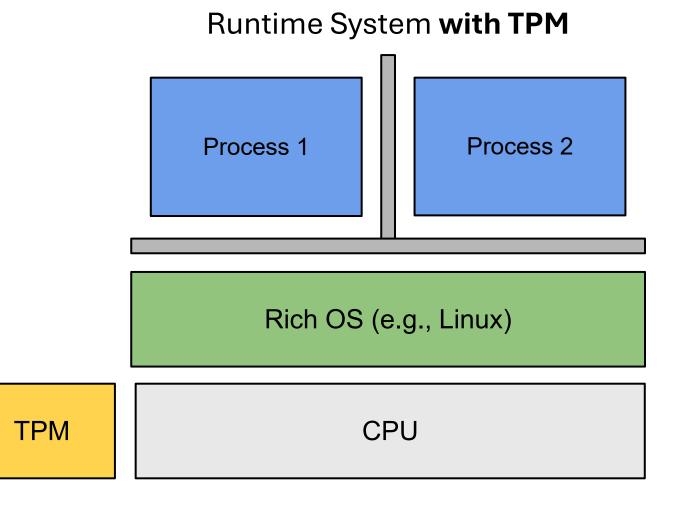
ARM Processors' TEE

- Splits the system into two worlds
- From ARM:
 - "The security of the system is achieved by partitioning all (...) hardware and software resources so that <u>they exist in one of two worlds</u> <u>the Secure world for the security subsystem</u>, and the <u>Normal world for everything else."</u>

Some visualizations....

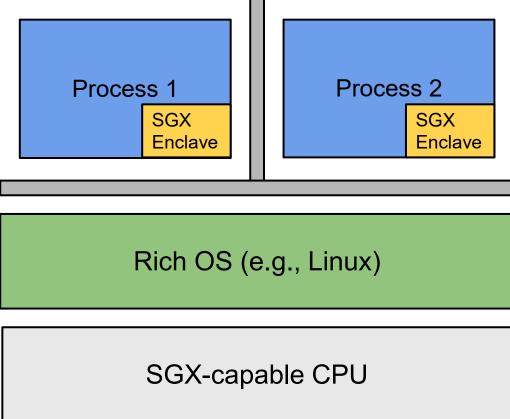


Some visualizations....

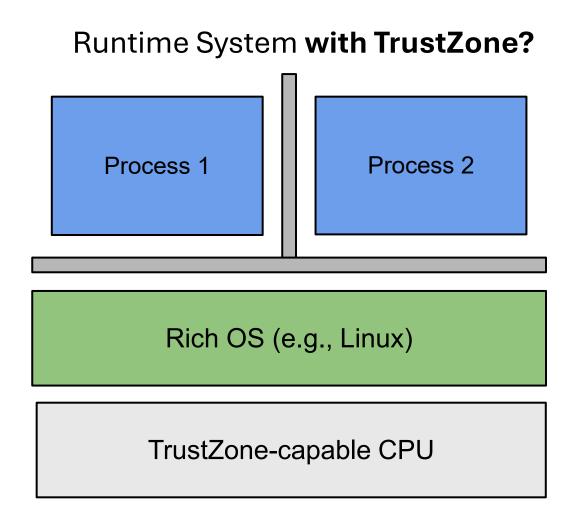


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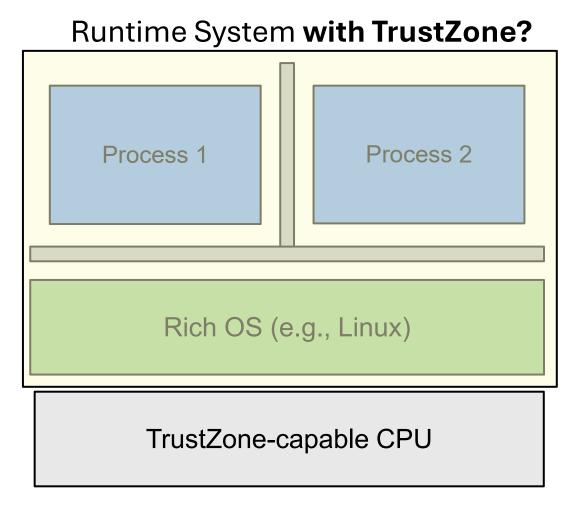
Runtime System with SGX



Some visualizations....



Some visualizations....



Split everything between two worlds

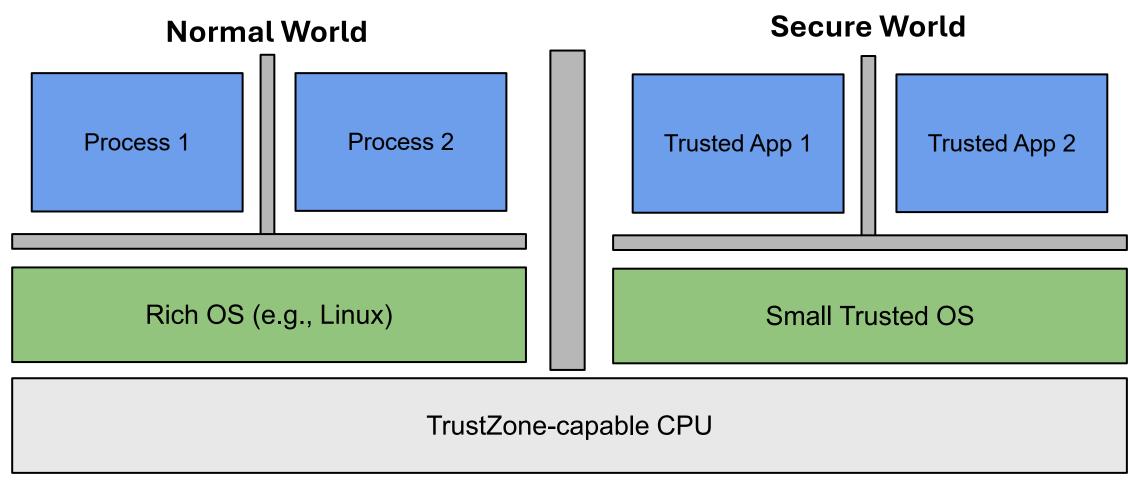
Some visualizations....

Runtime System with TrustZone?

Process 1 Process 2 Rich OS (e.g., Linux)

TrustZone-capable CPU

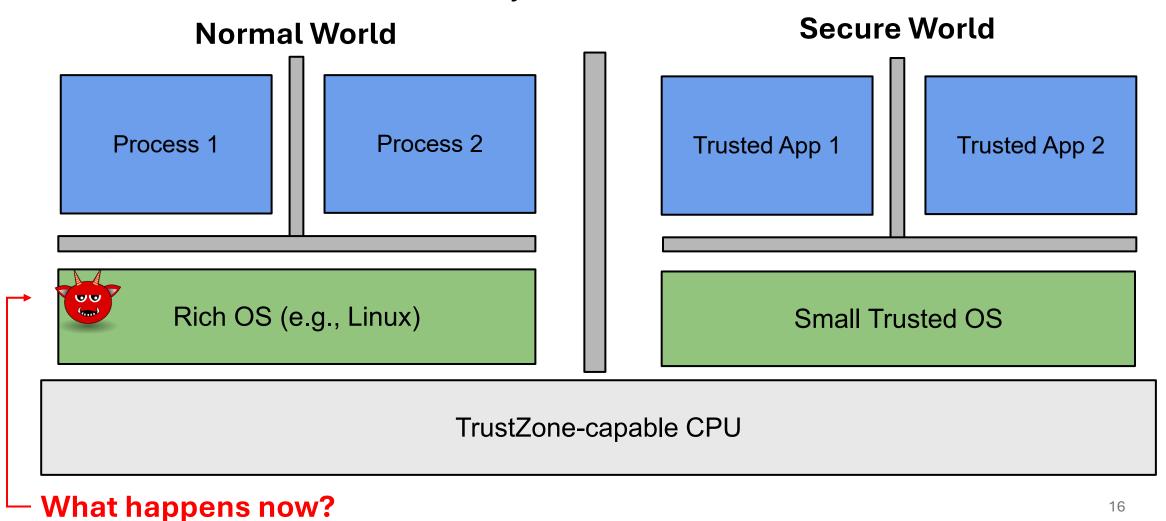
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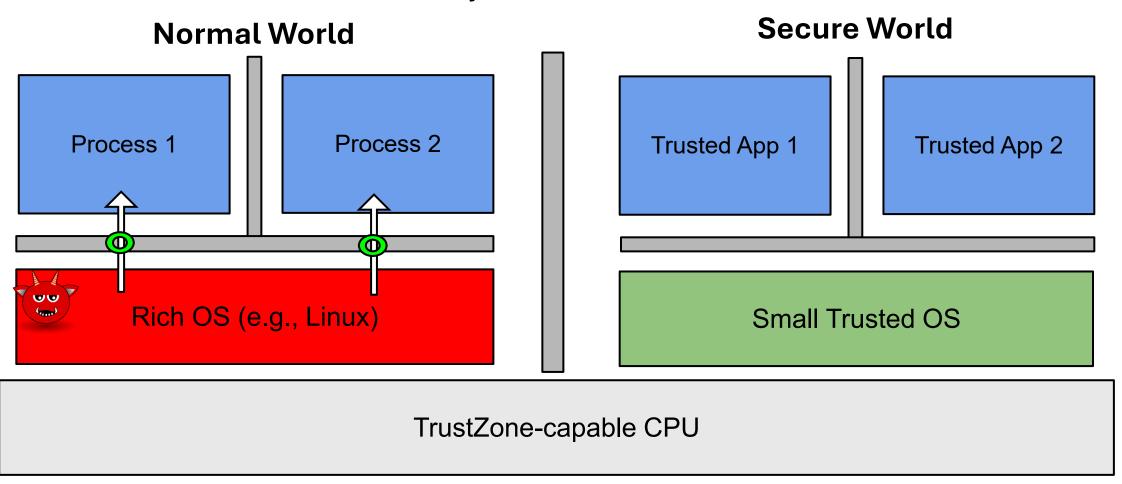
TrustZone's Guarantee:

Even if the Normal World is fully compromised, the Secure World remains safe, confidential, isolated, etc.

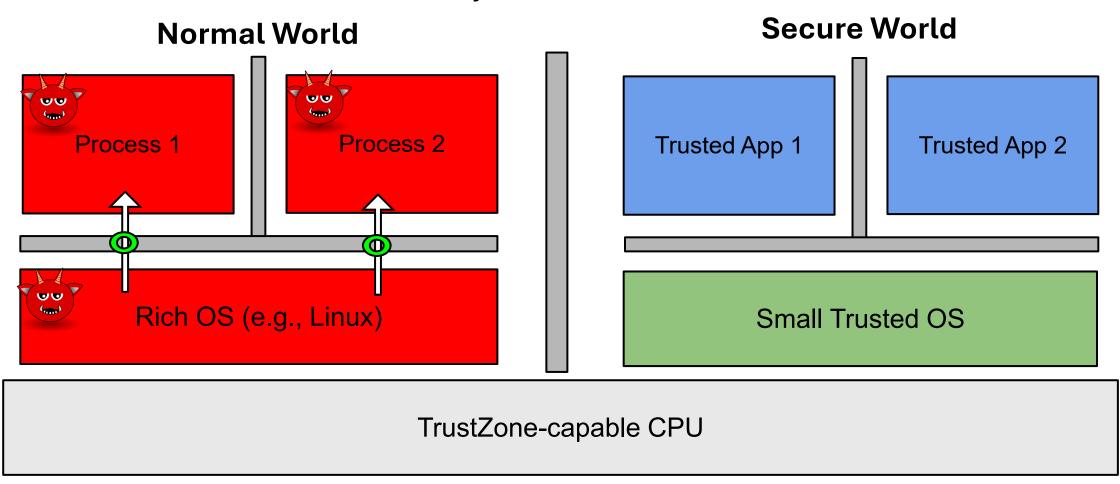
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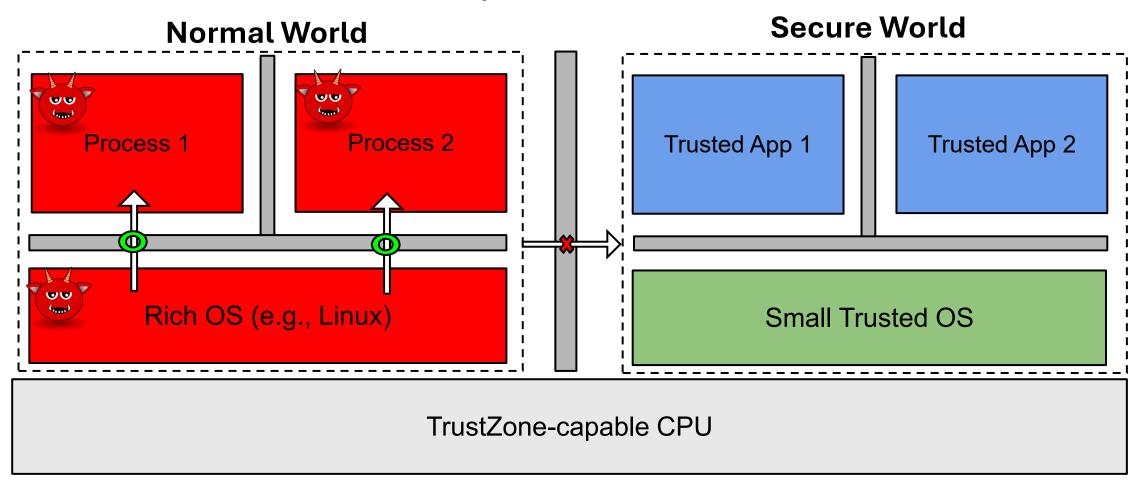
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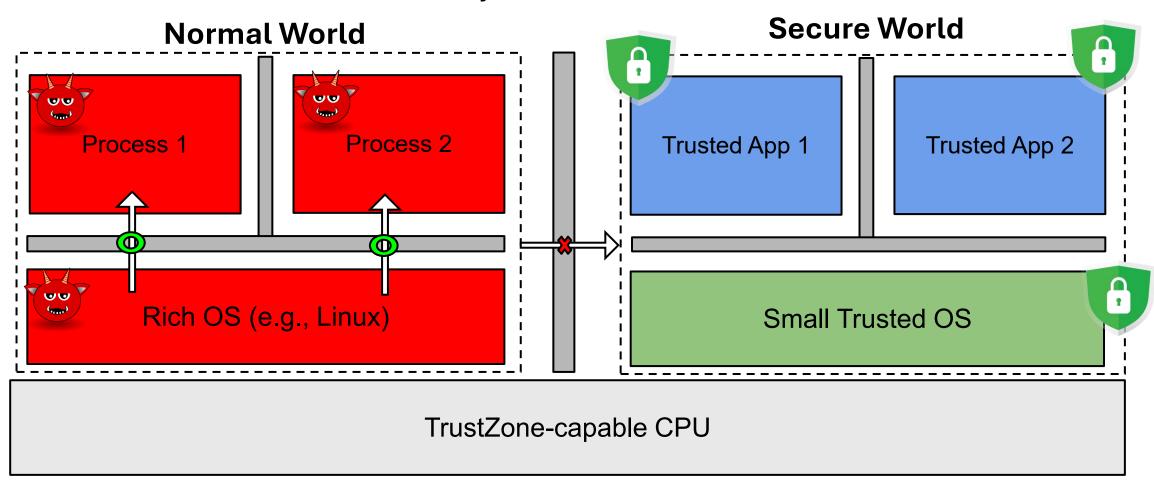
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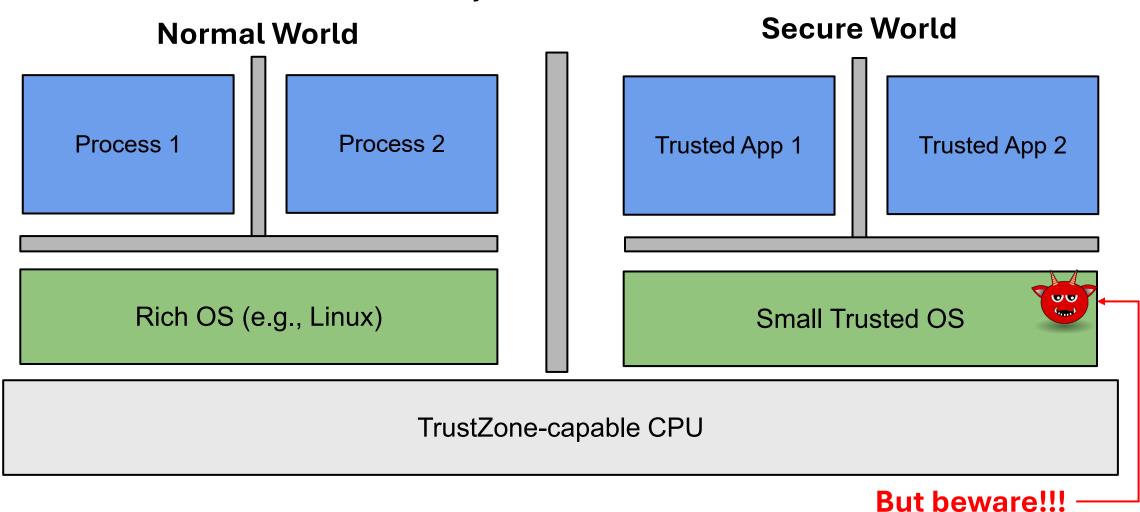
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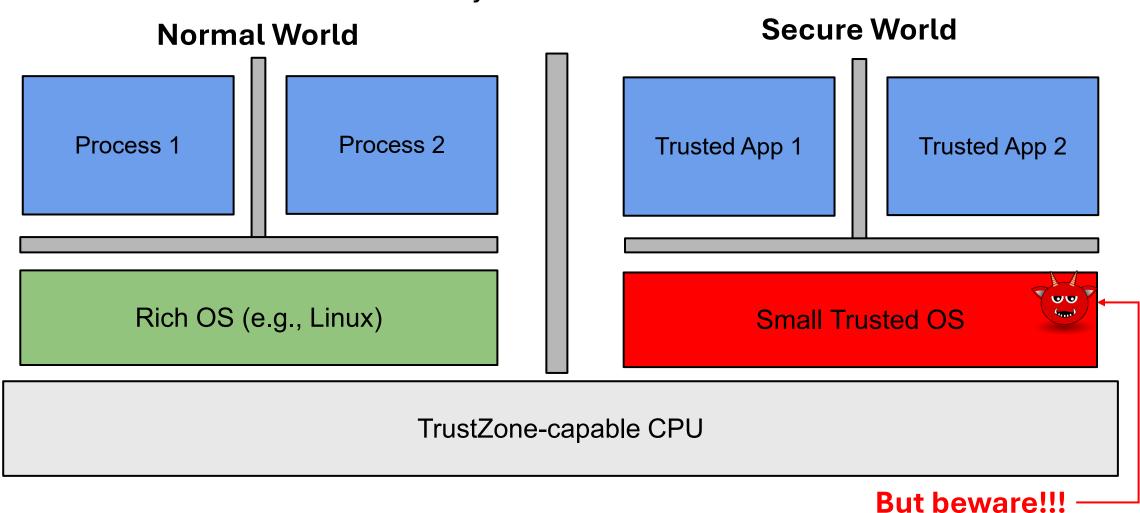


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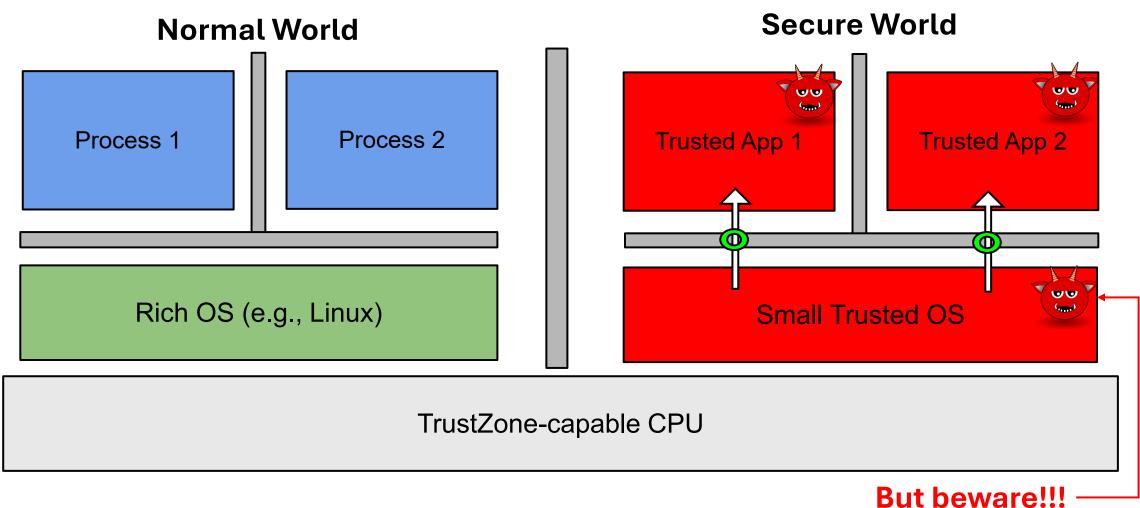


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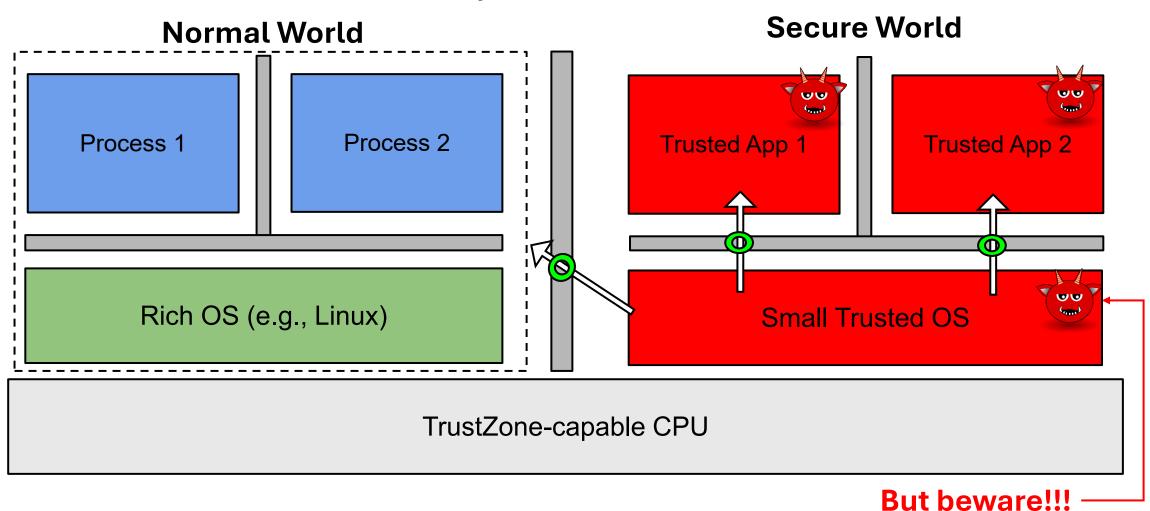




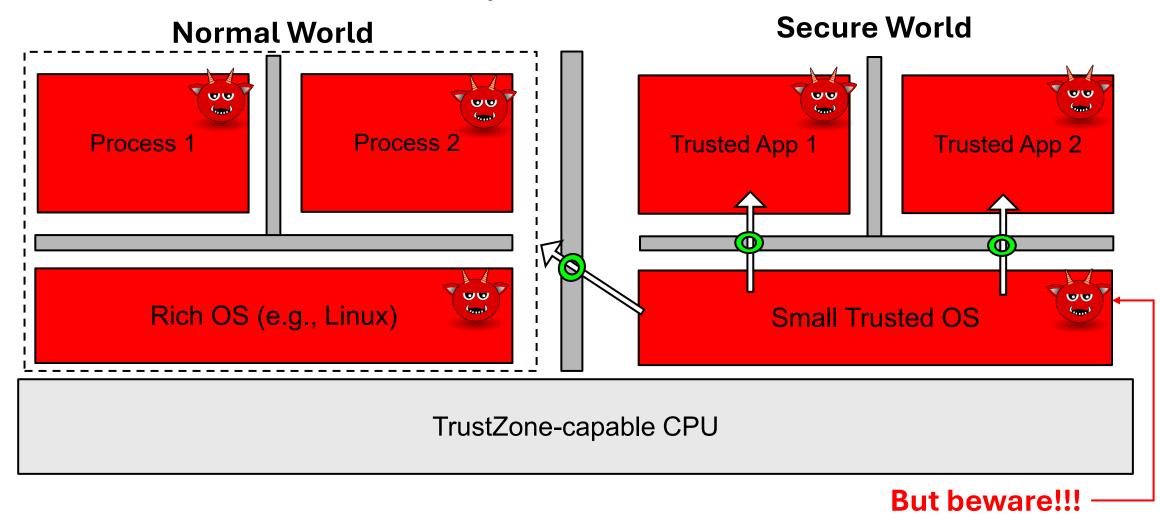
Some visualizations....



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Topics:

Isolation in TrustZone

Secure Monitor Calls (SMC) – Invocation of Secure World code

Android

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Isolation in TrustZone

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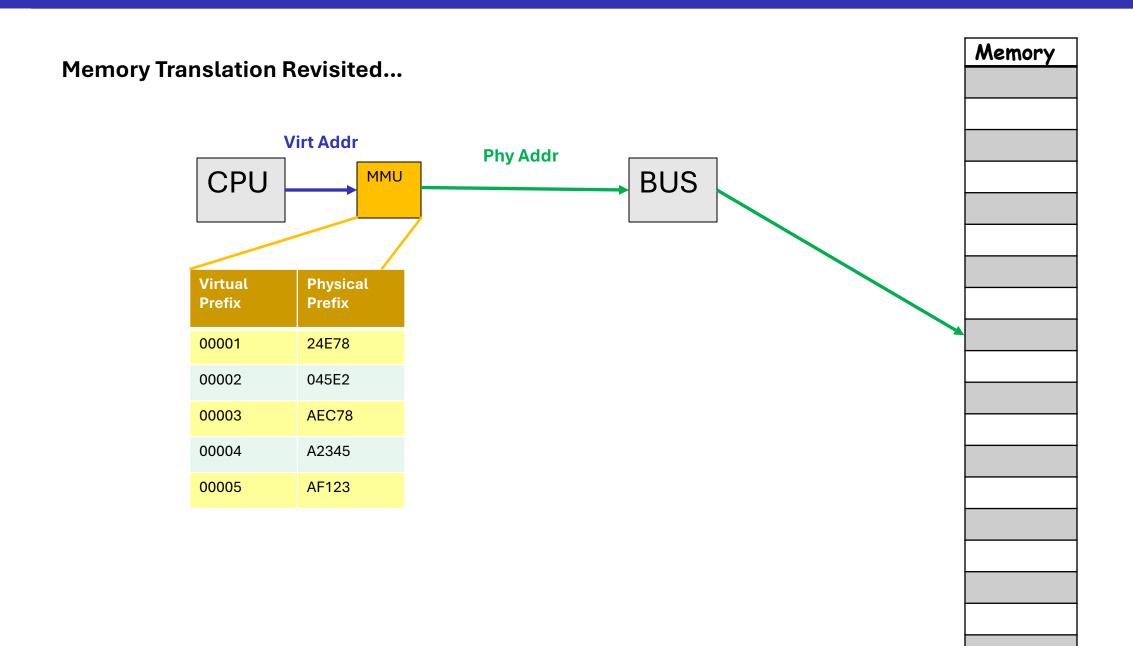
Android

- Store and manipulate security-critical info within the Secure World
 - Passwords, biometrics, private data, etc.

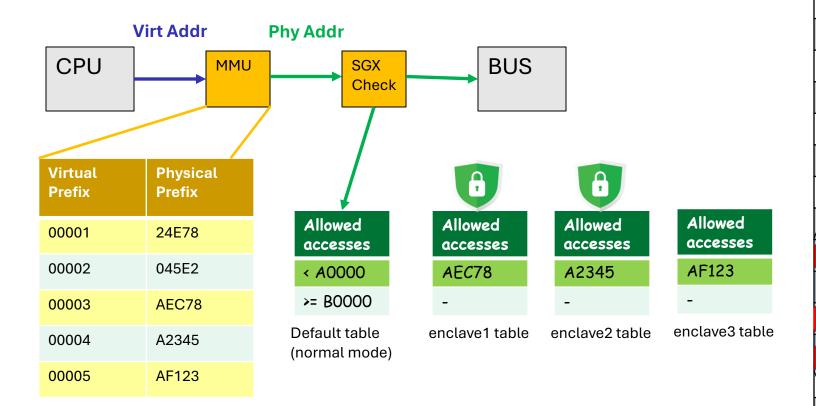
- Store and manipulate security-critical info within the Secure World
 - Passwords, biometrics, private data, etc.
- Keep the code inside the Secure World minimal \rightarrow small TCB

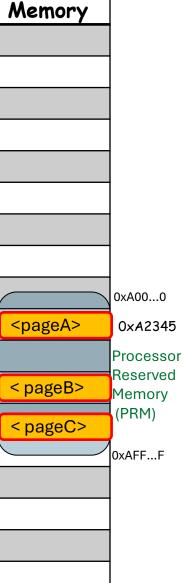
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 - Passwords, biometrics, private data, etc.
- Non-security tasks stay out in the Normal World
 - E.g., network stack, device drivers, UI implementation, etc.
- Normal World Apps make requests to Secure World apps via well-defined APIs
 - E.g., request decryption, check this biometric input, etc...

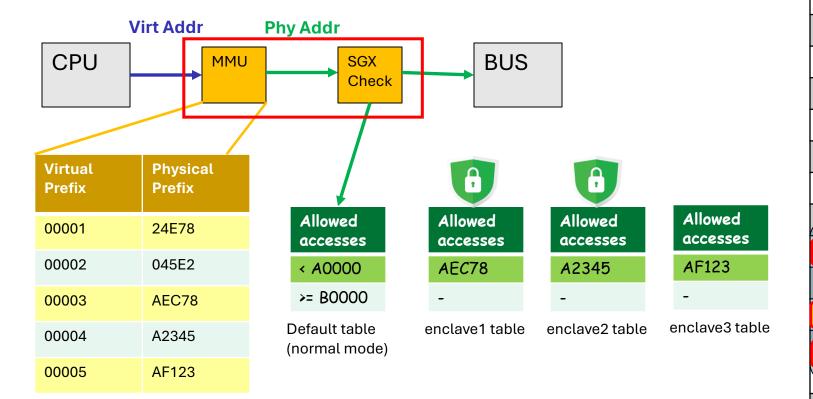


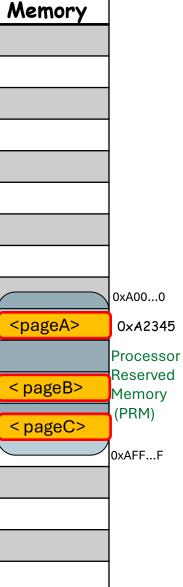
Memory Translation in Intel SGX Revisited....





Memory Translation in Intel SGX Revisited....





Similar idea -> CPU is now involved in the memory translation

TrustZone approach:

- Two worlds \rightarrow two page tables
- Both are active in an MMU at a given time

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- One additional bit in the CPU

 tells MMU which table to load

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TrustZone approach:

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 - Non-Secure (NS) bit:

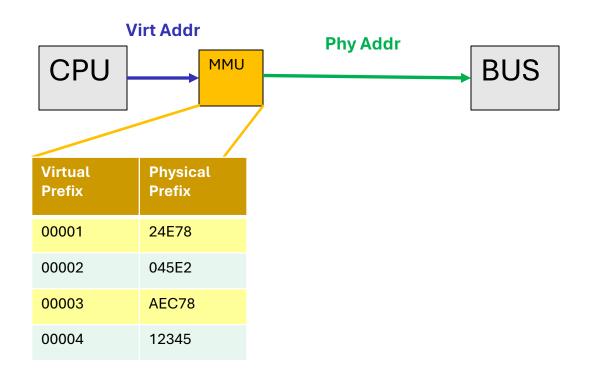
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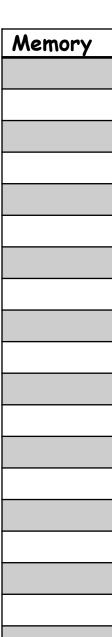
TrustZone approach (part 1):

- Two worlds \rightarrow two page tables
- Both are active in an MMU at a given time
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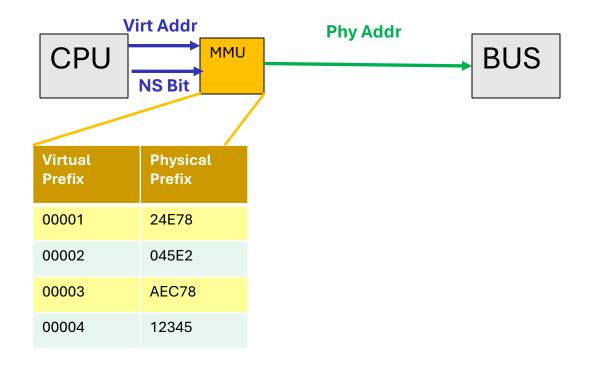
 tells MMU which table to load
 - Non-Secure (NS) bit:
 - NS = 1 → currently in Normal World, Secure World access is blocked
 - NS = 0 → currently in Secure World, Secure World access is allowed

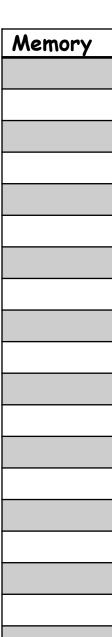
The details...



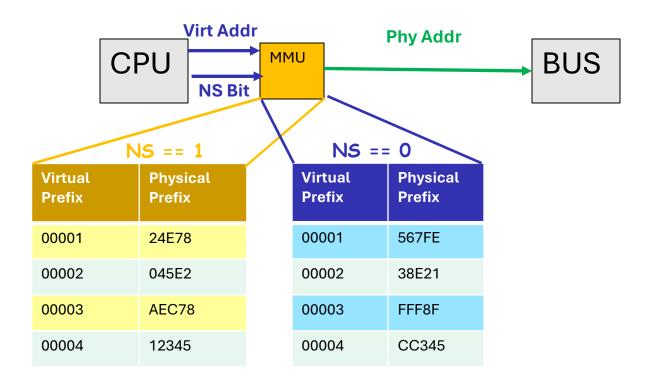


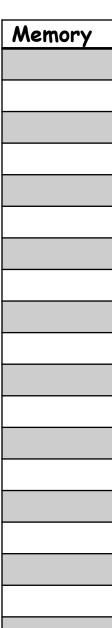
Now, the CPU also passes the NS bit to the MMU



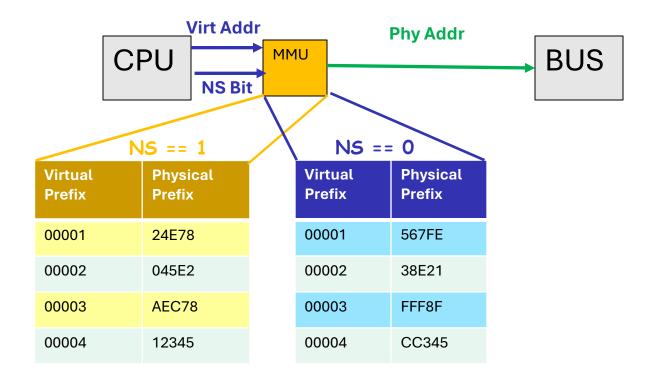


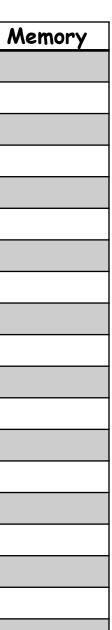
And the MMU has two page tables. NS bit tells MMU which to use





Now, the CPU also passes the NS bit to the MMU

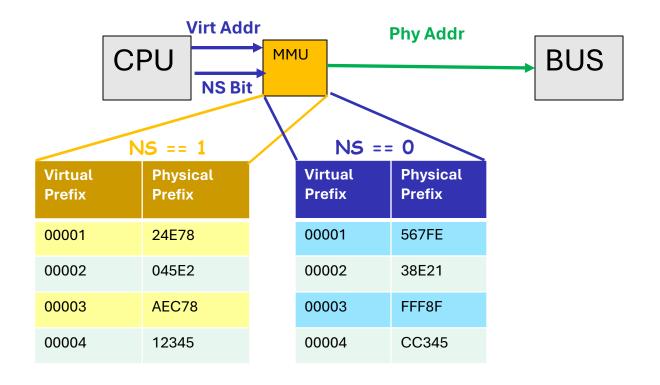


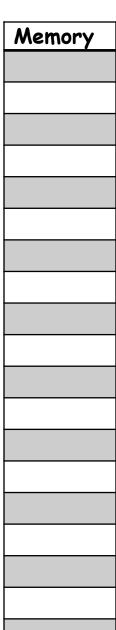


TrustZone approach (continued..)

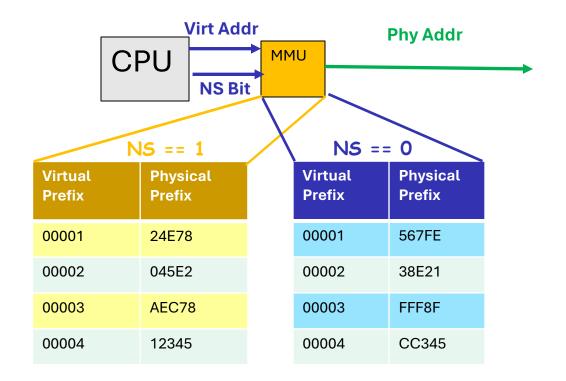
- Physical Memory Partitioning in addition to the modified MMU!
- TrustZone enables configuration of specific <u>physical</u> memory regions as <u>secure</u> or <u>non-secure</u>, such that applications can only access memory assigned to their world
- How?
 - Two hardware controllers:
 - TrustZone Address Space Controller (TZASC) → on chip memory (SoC) and DRAM
 - TrustZone Memory Adapter (TZMA) → off-chip memory (e.g., external peripherals SRAM)
- TZASC and TZMA have the same function applied to different resources

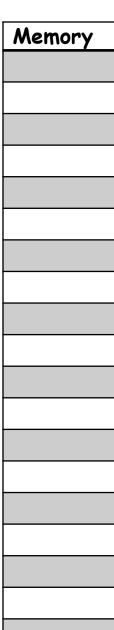
Where are the TZASC/TZMA?

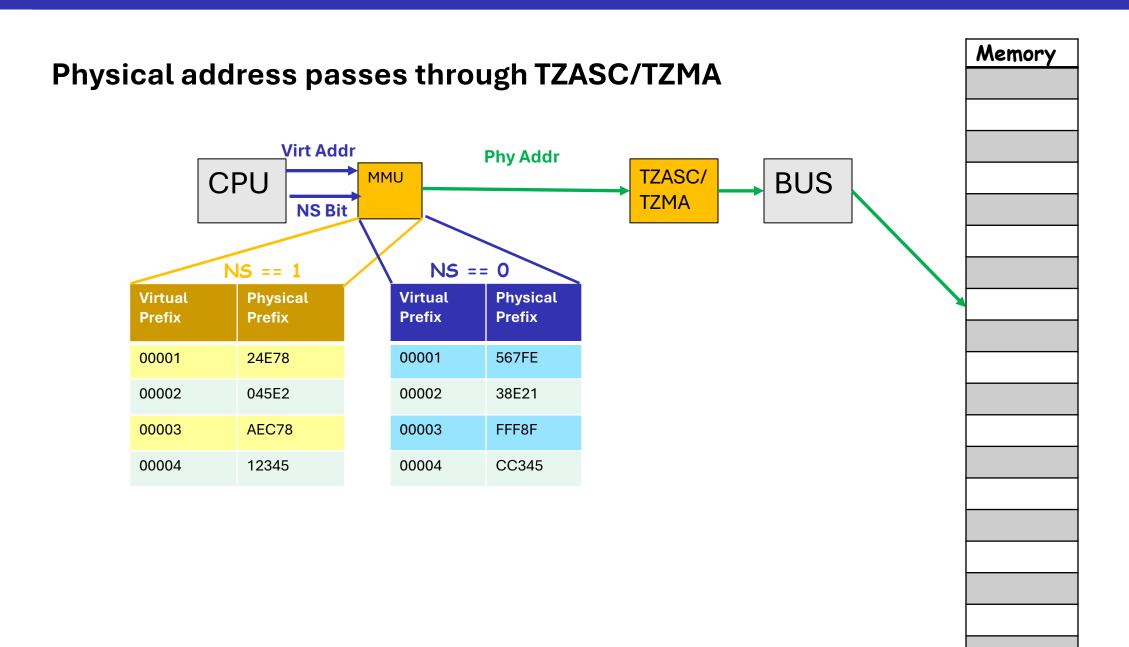


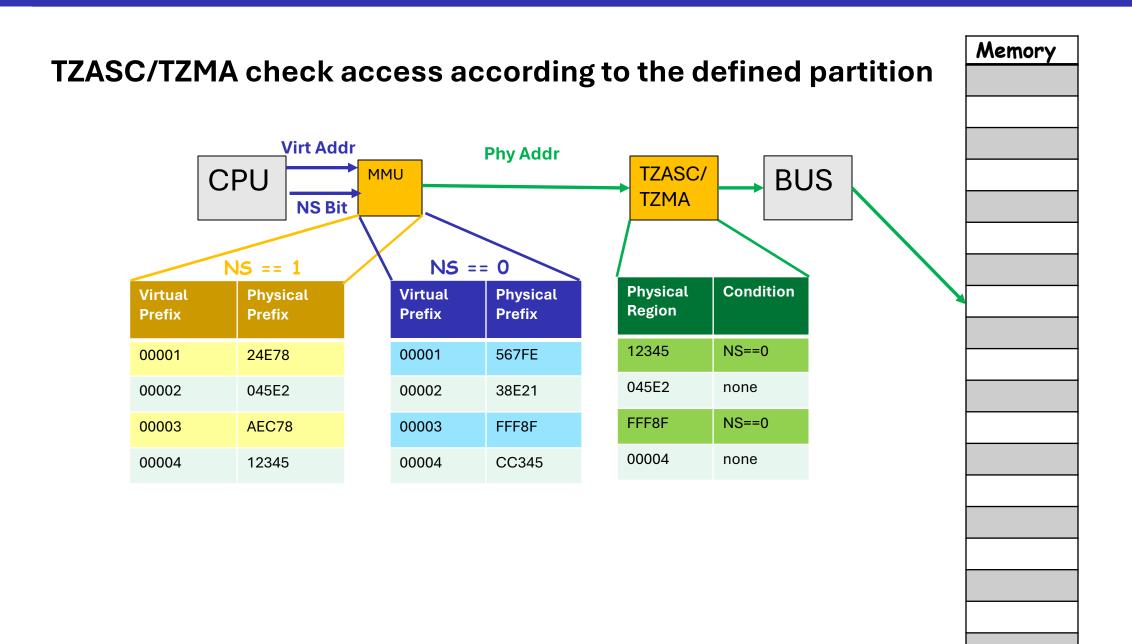


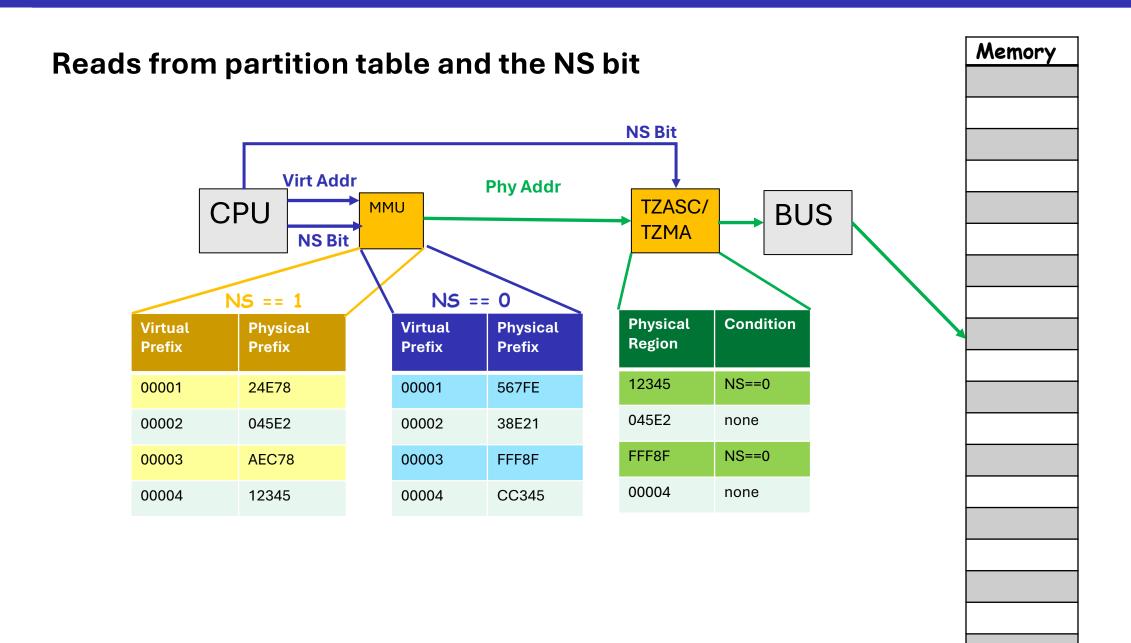
Instead of the MMU accessing the BUS directly...

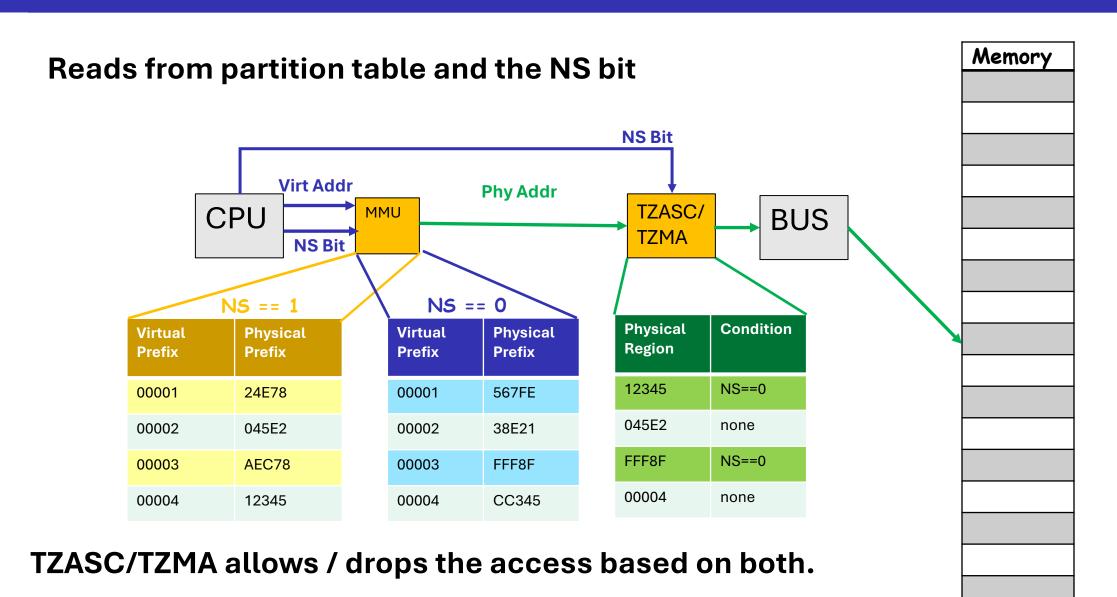












TrustZone-A MMU + TZASC/TZMA

Together provide isolation in TrustZone

TZASC/TZMA:

Implement physical isolation between the worlds

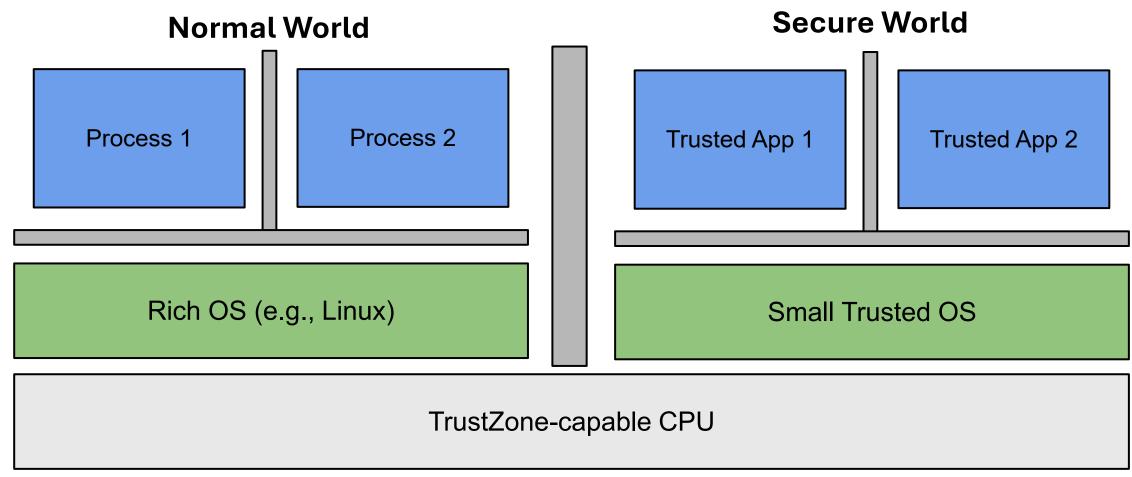
- Isolate physical memory, peripherals, and hardware resources
- Provides system-level isolation

MMU:

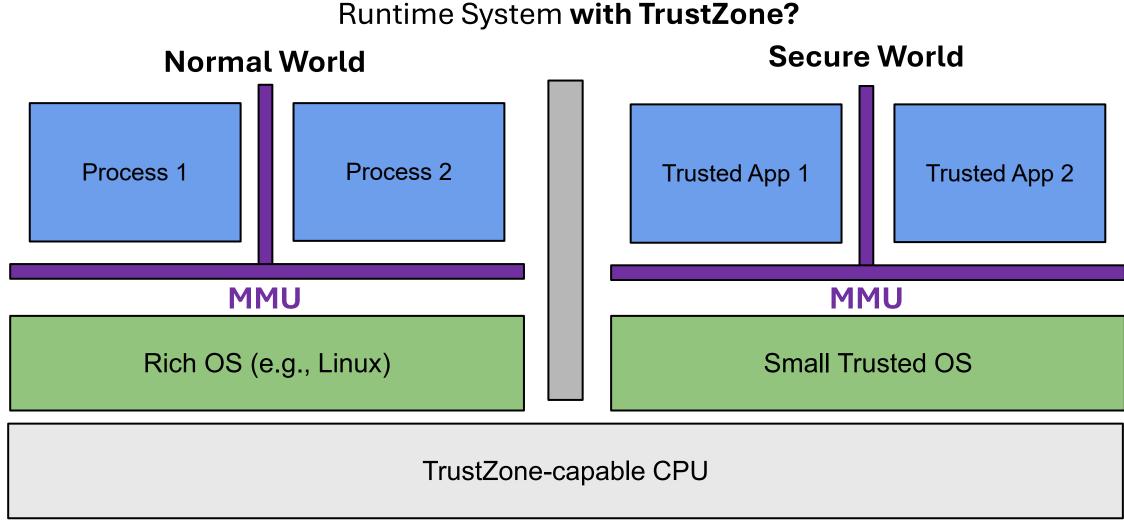
- Virtual isolation between processes running in each world
- Provides process-level isolation

Some visualizations....

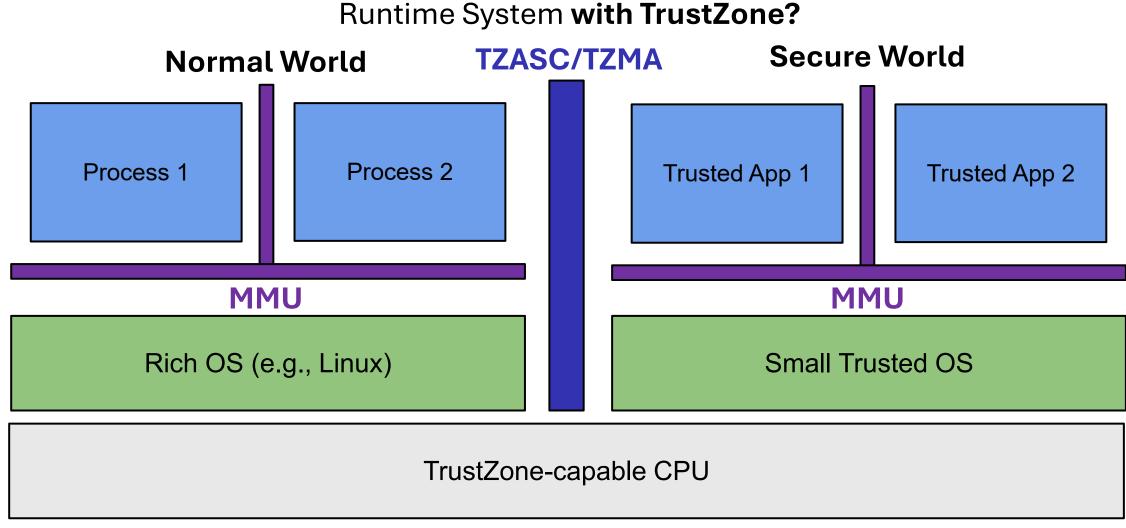
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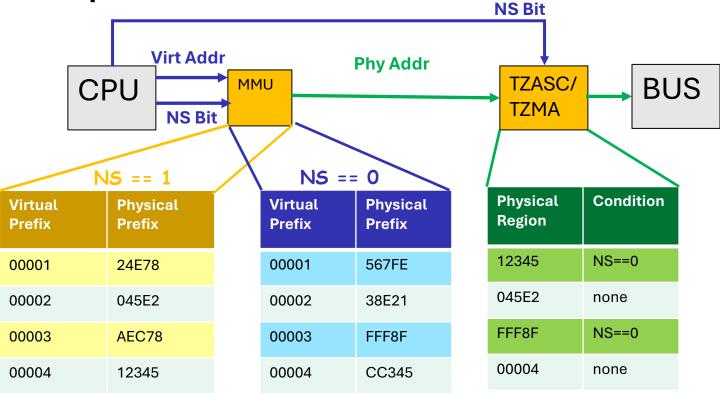
Some visualizations....



Some visualizations....



Some examples....



Access 1: NS = 1; Virt Addr = 0x00002123

Access 2: NS = 0; Virt Addr = 0x00003456

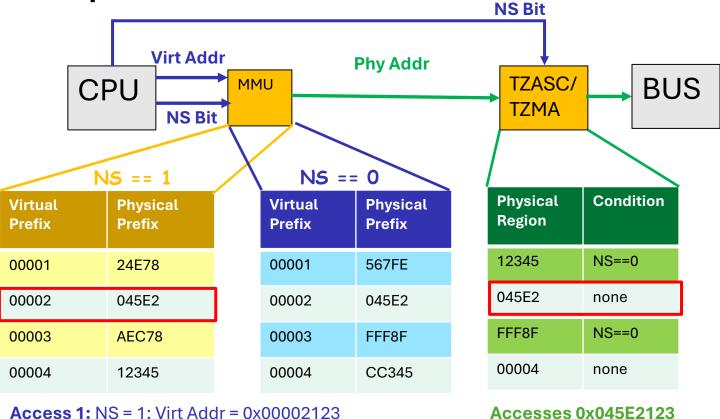
Access 3: NS = 1; Virt Addr = 0x00004789

Access 4: NS = 1; Virt Addr = 0x00006333

Access 5: NS = 0; Virt Addr = 0x00002123

Memory

Some examples....



Access 1: NS = 1; Virt Addr = 0x00002123

Access 2: NS = 0; Virt Addr = 0x00003456

Access 3: NS = 1; Virt Addr = 0x00004789

Access 4: NS = 1; Virt Addr = 0x00006333

Access 5: NS = 0; Virt Addr = 0x00002123

Memory

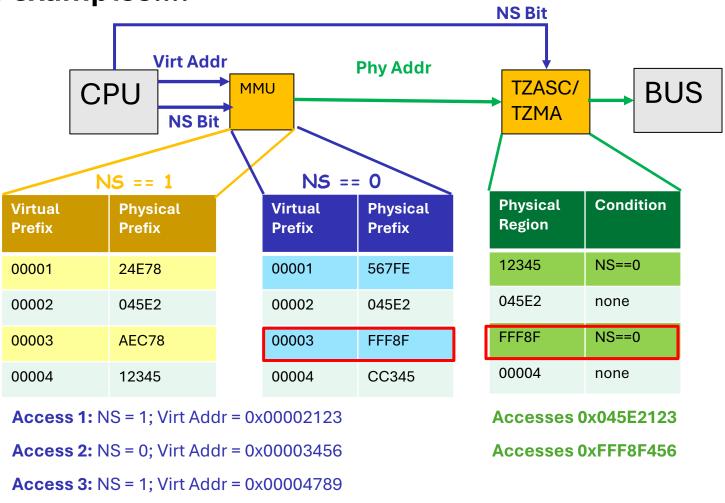
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ARM TrustZone – Memory Translation

Some examples....

Access 4: NS = 1; Virt Addr = 0x00006333

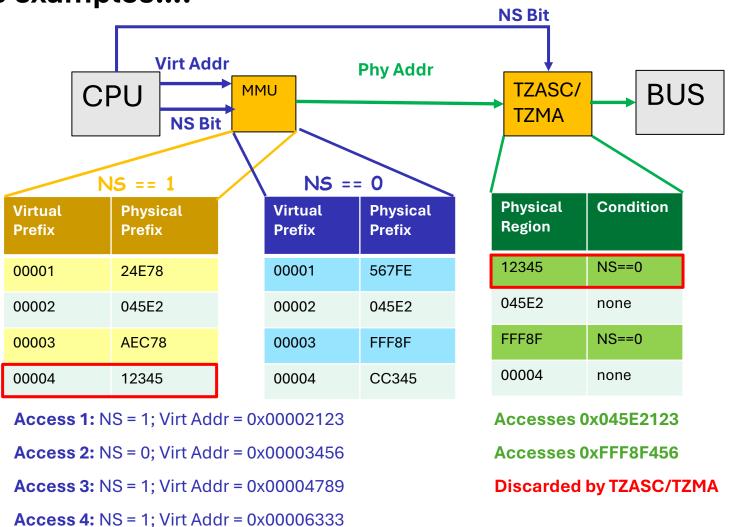
Access 5: NS = 0; Virt Addr = 0x00002123



Memory

Access 5: NS = 0; Virt Addr = 0x00002123

Some examples....

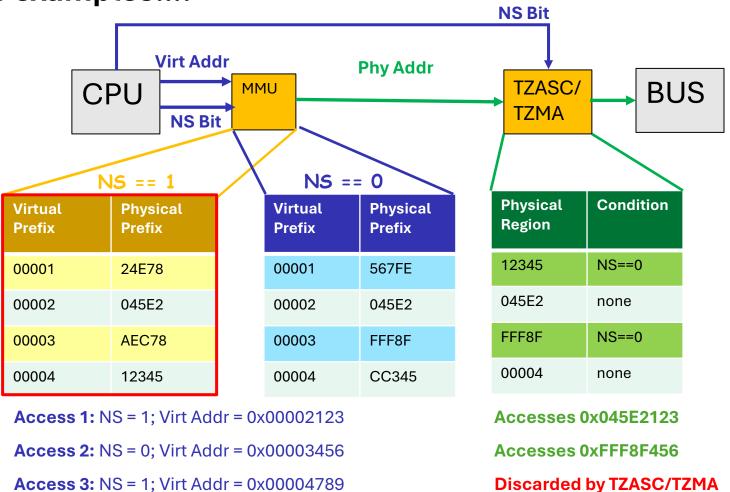


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Access 4: NS = 1; Virt Addr = 0x00006333

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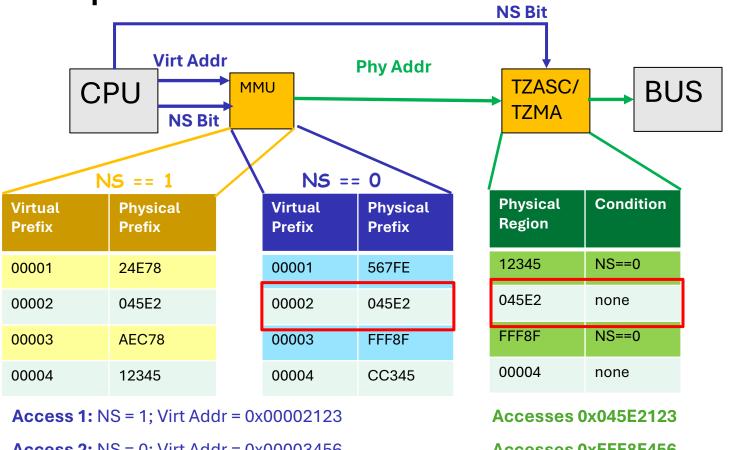
Some examples....



Discarded by MMU

Memory	

Some examples....



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Access 4: NS = 1; Virt Addr = 0x00006333

Access 5: NS = 0; Virt Addr = 0x00002123

Accesses 0xFFF8F456

Discarded by TZASC/TZMA

Discarded by MMU

Accesses 0x045E2123

Memory

Two key questions:

Who configures the TZASC/TMA table?

• Who controls the NS bit value?

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- Who configures the TZASC/TMA table?
 - Secure World code: first configuration after boot! It is more privileged than Normal World
 - Secure world executes first and configures TZMA/TZASC before launching the normal world and rich OS.
 - Security of TrustZone requires TrustZone-aware Secure boot!
- Who controls the NS bit value?

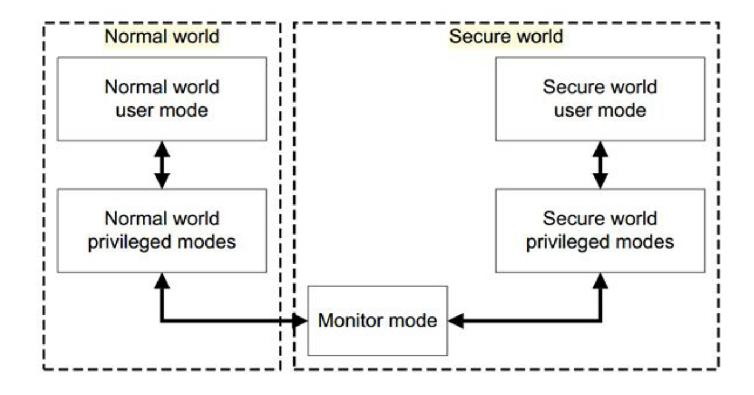
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- Who controls the NS bit value?
 - The CPU in hardware
 - From normal world, NS bit can only be changed (1 \rightarrow 0) by issuing a **Security Monitor Call (SMC)**
 - SMC atomically gives control to Secure World and sets NS=0
 - SMC jumps to **Security Monitor** that performs context switch between the worlds
 - The NS bit is set back to NS=1 before returning to Normal World

Security Monitor and SMC:

Switching between worlds requires a security monitor call (SMC)

The Security Monitor is part of the Secure World's TCB



Caching in TrustZone

The problem: the CPU, and consequently the cache, must be securely shared between worlds

The TZMA/TZASC split physical memory, but not the cache

Caching in TrustZone

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The TZMA/TZASC split physical memory, but not the cache

So without any additional measures, the following is a possibility:

- 1. Secure World is running
- 2. Secure World transfers context back to the Normal World
- 3. Normal World reads the same cached address used by Secure World
- Data leaked! → Isolation is broken!

Caching in TrustZone

How to handle this problem?

Naïve solution: Remove the cache => secure, but extremely slow.

<u>Alternative:</u> Always flush the cache when switching worlds => secure, but still pretty slow.

Caching in TrustZone

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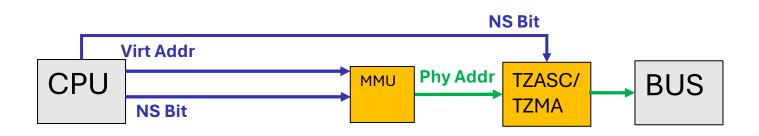
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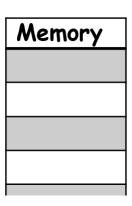
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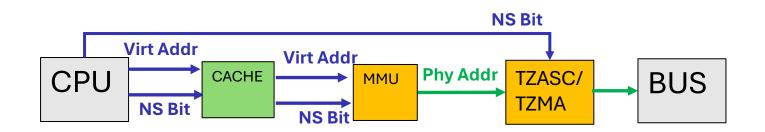
TrustZone's solution:

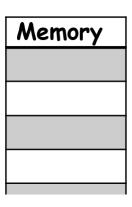
Include the NS bit in the cache look-up => no need to flush the cache!

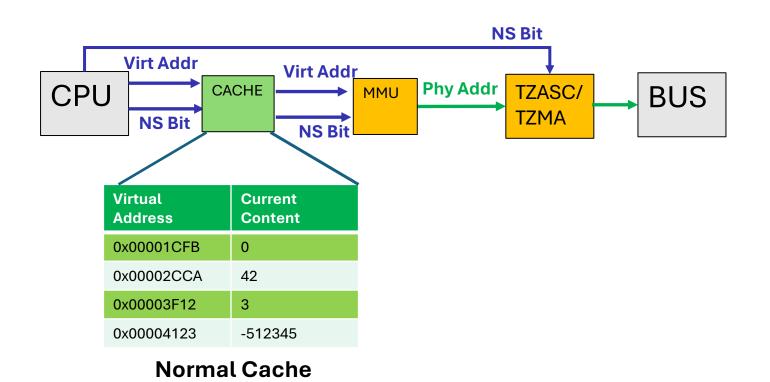
- Allows for fast world switching
- Cached data may be kept across successive switches

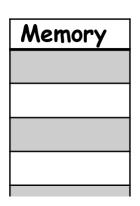


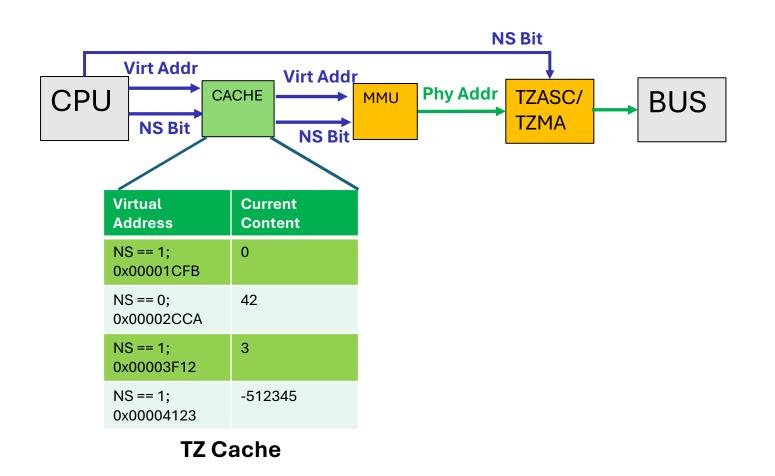


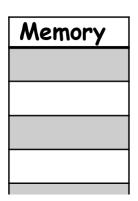




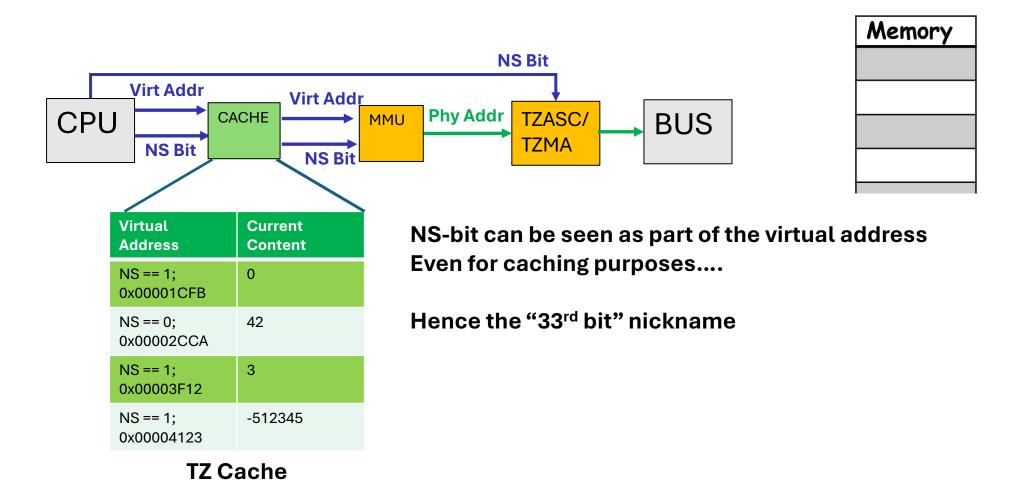


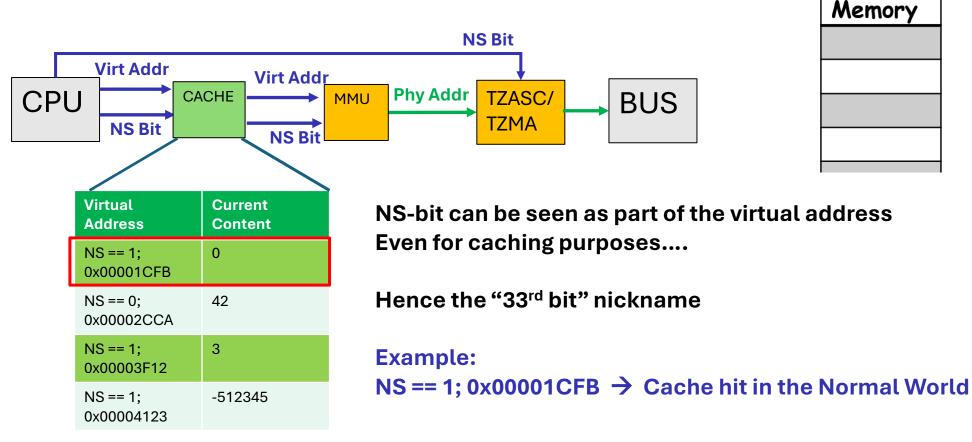




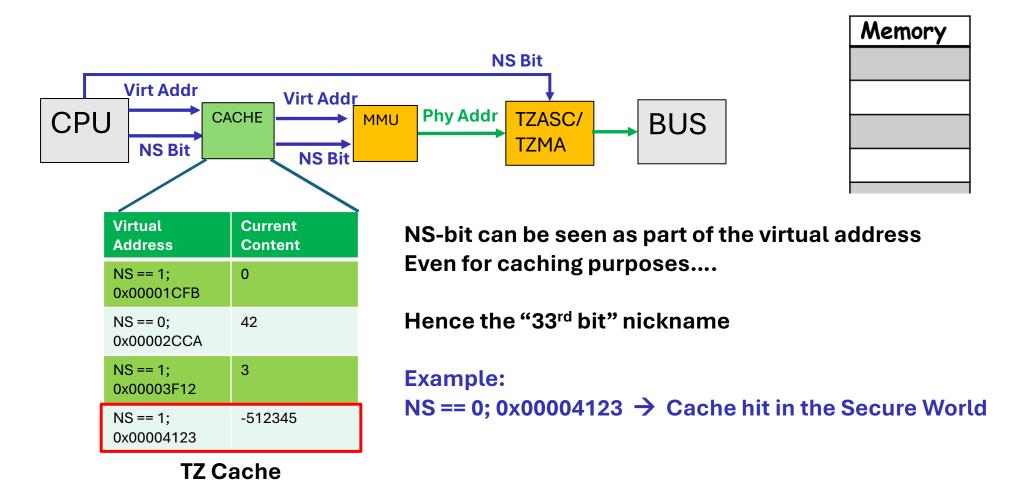


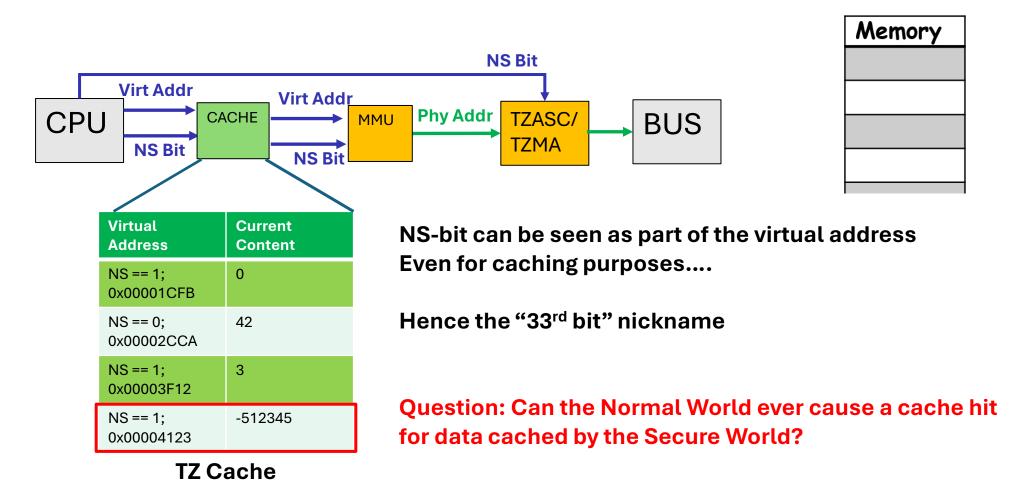
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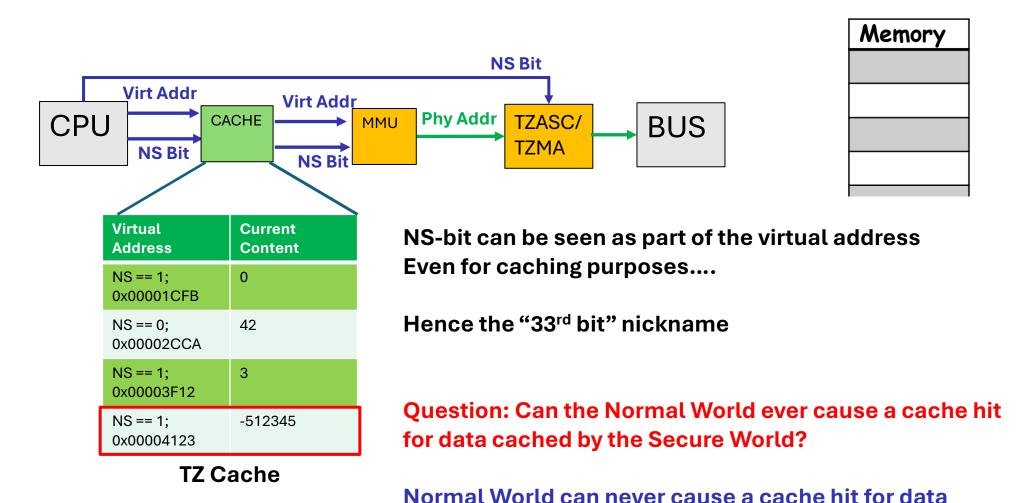


TZ Cache





Caching in TrustZone



cached by the Secure World execution.

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Caching in TrustZone

Also, the MMU has a cache:

- Called the TLB: Translation Lookaside Buffer
- Same principle for TrustZone's CPU cache

Summary of Isolation in TrustZone:

In main memory:

Physical isolation implemented by TZMA/TZASC based on NS-bit

Within CPU cache and TLB:

- Propagate NS-bit through every virtual address look-up
- NS-bit is "33rd bit"

Reminder:

• NS bit value is controlled by CPU hardware. Only way to set it to 0 is by calling SMC, which also gives control to TrustZone's trusted Security Monitor

ARM TrustZone Overview

Topics:

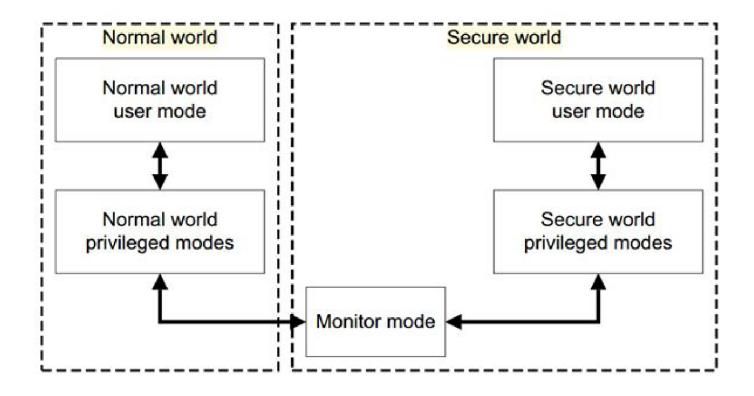
Isolation in TrustZone

Secure Monitor Calls (SMC) – Invocation of Secure World code

Android

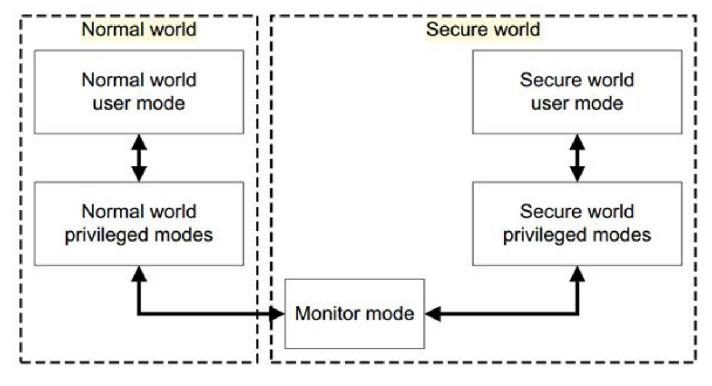
Revisiting the system flow:

Controlled enter and exit from the Secure World



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Controlled enter and exit from the Secure World



The moment the CPU Flips the NS-bit

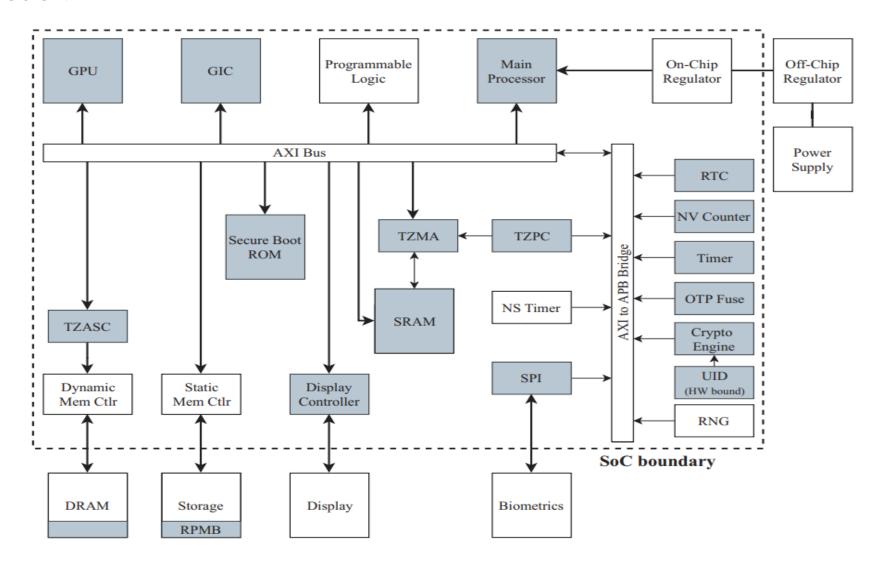
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 - Inputs destined to the secure world

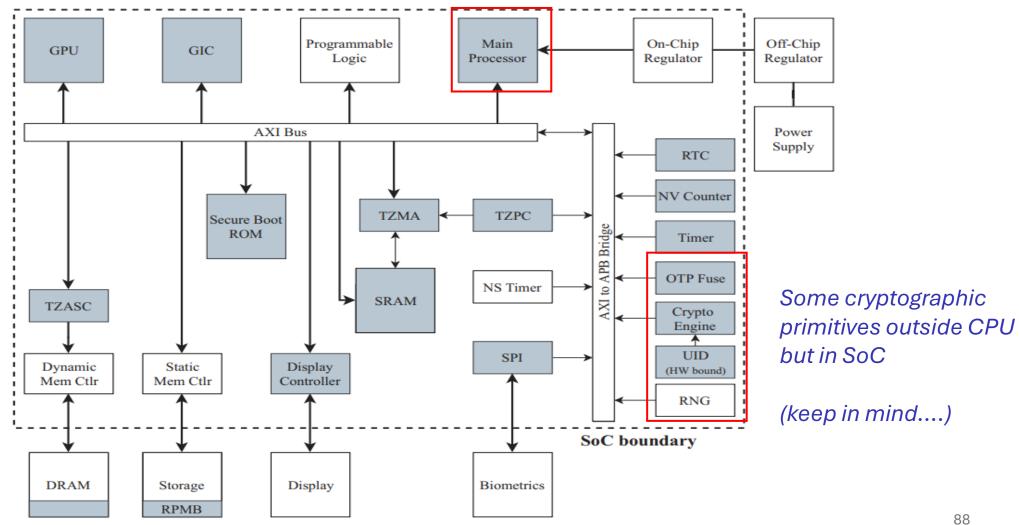
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- SMC Arguments are passed in registers R1-R7
 - Inputs destined to the secure world
- Results are also returned to normal world using registers R0-R7

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 - Can be used to tell the security monitor which Trusted App is the destination of this call
- SMC Arguments are passed in registers R1-R7
 - Inputs destined to the secure world
- Results are also returned to normal world using registers R0-R7
- Convention: not enforced by hardware anywhere
 - It is up to the Security Monitor to define its own behavior
 - Must then be followed/implemented by the SMC caller

The whole beast:



The whole beast:



Important reminders about TrustZone's design:

- Secure boot must guarantee that the Secure World runs first
 - After Secure World completes secure boot → "ACTIVE"
- Availability
 - Boots first!
 - Also, resources assigned to secure world have priority (e.g., interrupts via TrustZone's GIC)
 - Different from TPM and SGX -> Has an "active" characteristic

ARM TrustZone Overview

Topics:

Isolation in TrustZone

• Secure Monitor Calls (SMC) – Invocation of Secure World code

Android

Provides runtime environment built atop TrustZone

- Android OS & Apps → in the normal world
- Trusted OS & Trusted Apps -> in the secure world

Features of interest:

Key store

Android Key Store:

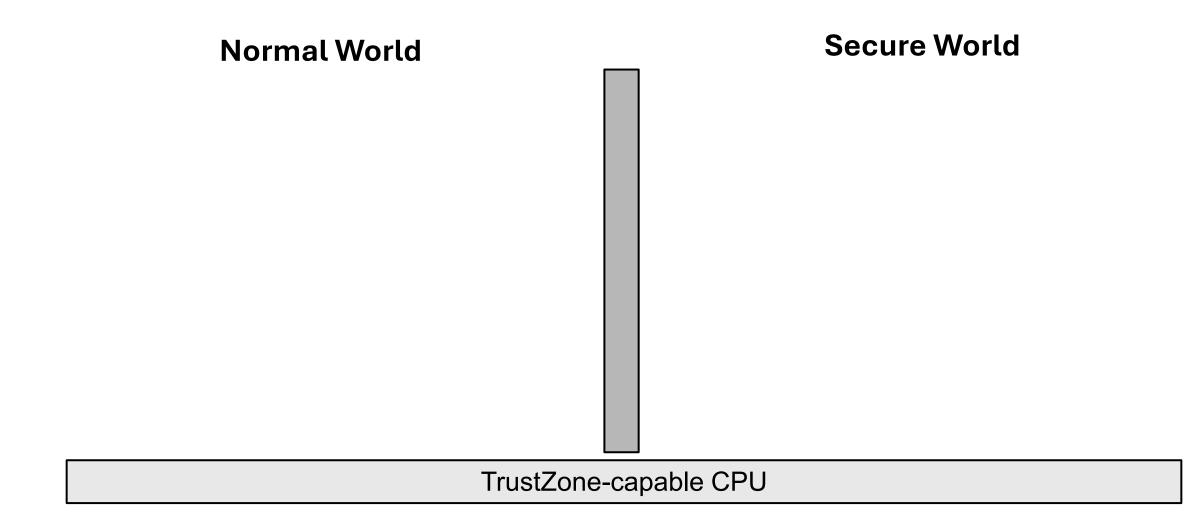
Protects key material from unauthorized use in two ways. First it .. Prevents the extraction of keys from application processes and from the Android device, and Second it makes apps specify the authorized use of their keys within the device and enforces those restrictions outside of the app's processes.

Android Key Store:

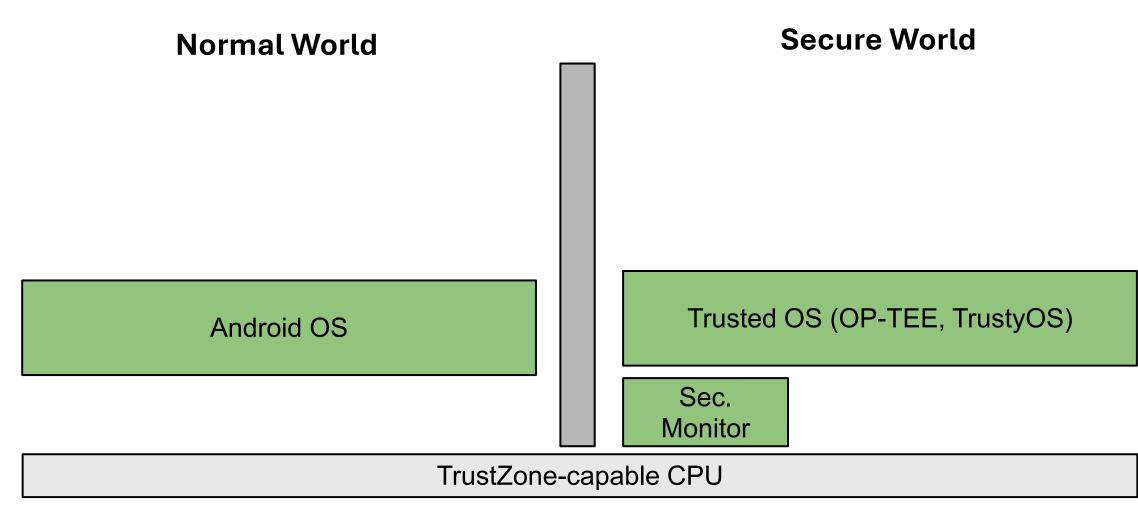
Extraction prevention is provided based on two security measures:

- Key material never enters the application process
 - Inputs for a operation that requires the key are fed into a "secure process"
 - Compromised App can use keys, but cannot extract the key itself
 - Confidentiality
- Keys can be bound to the TEE
 - Similar to "wrap key" to be used from a particular device
 - Integrity

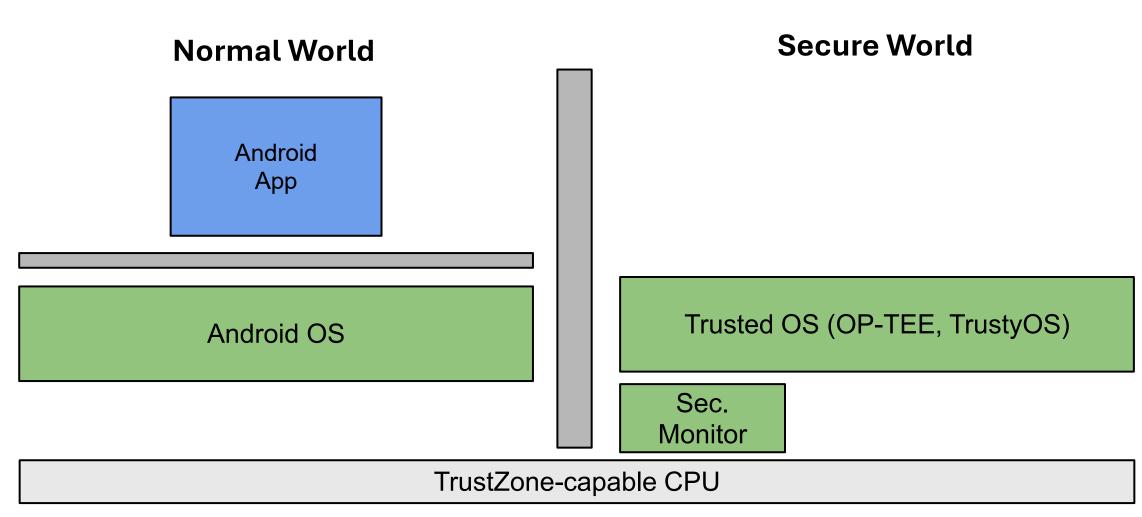
How does Key store work? First, lets setup the key players...



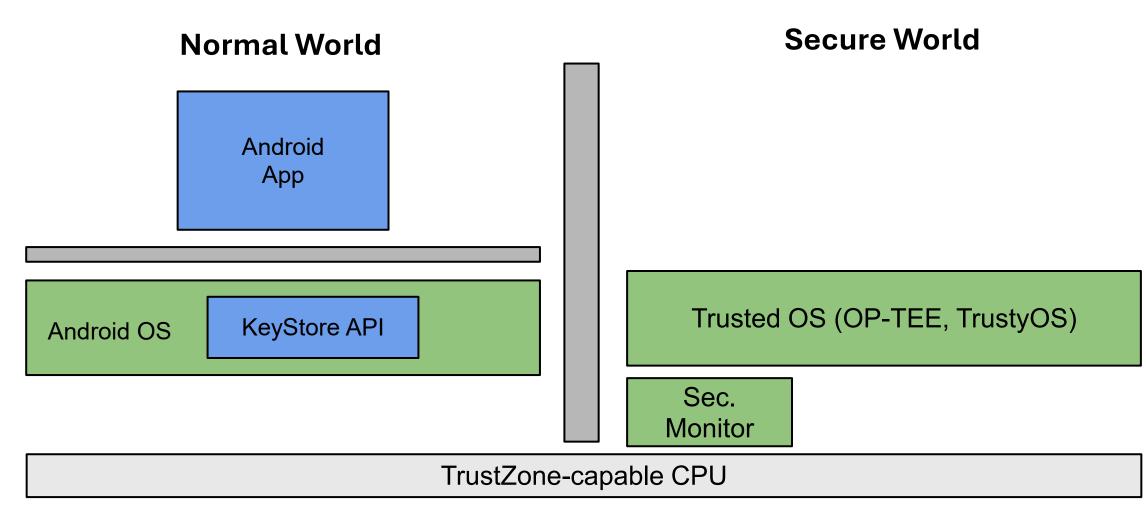
First OSes: Android OS in the Normal World, and a Trusted OS in the Secure World



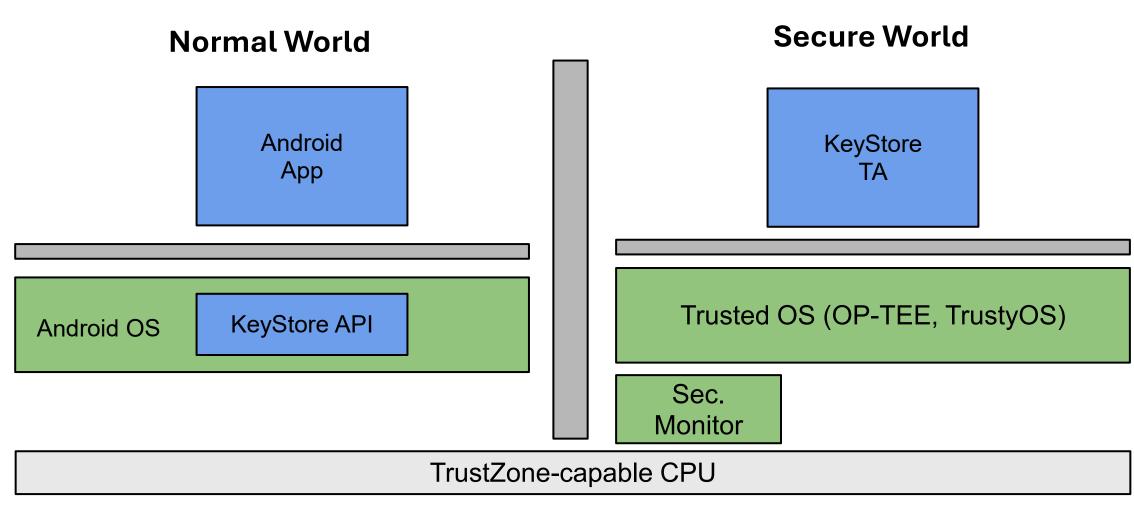
To simplify things, lets assume there is one Android app running in Normal World



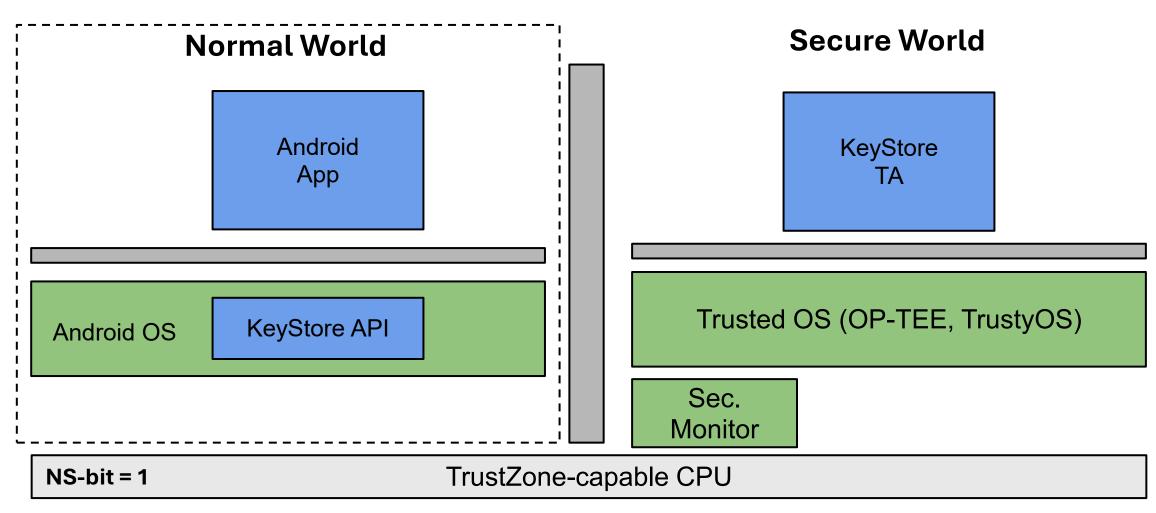
Within the Android OS is the KeyStore API



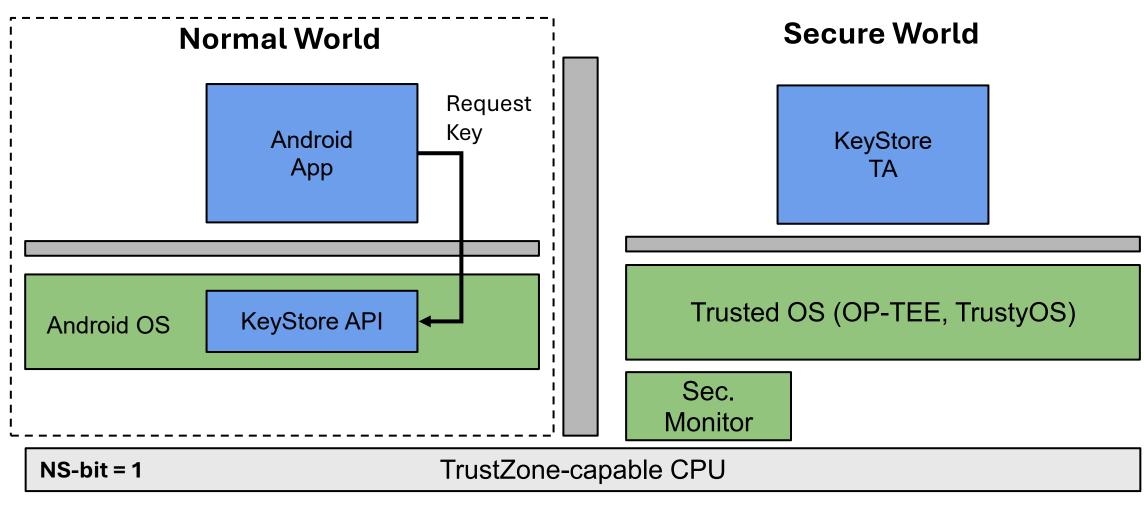
Within the Android OS is the KeyStore API, and a corresponding Keystore TA



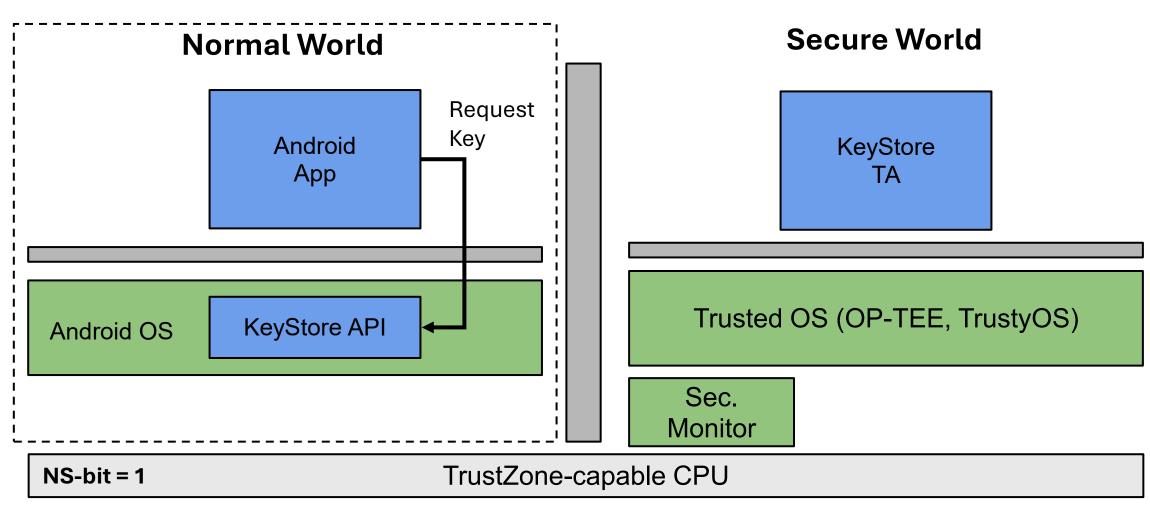
Lets assume the Android App is currently executing



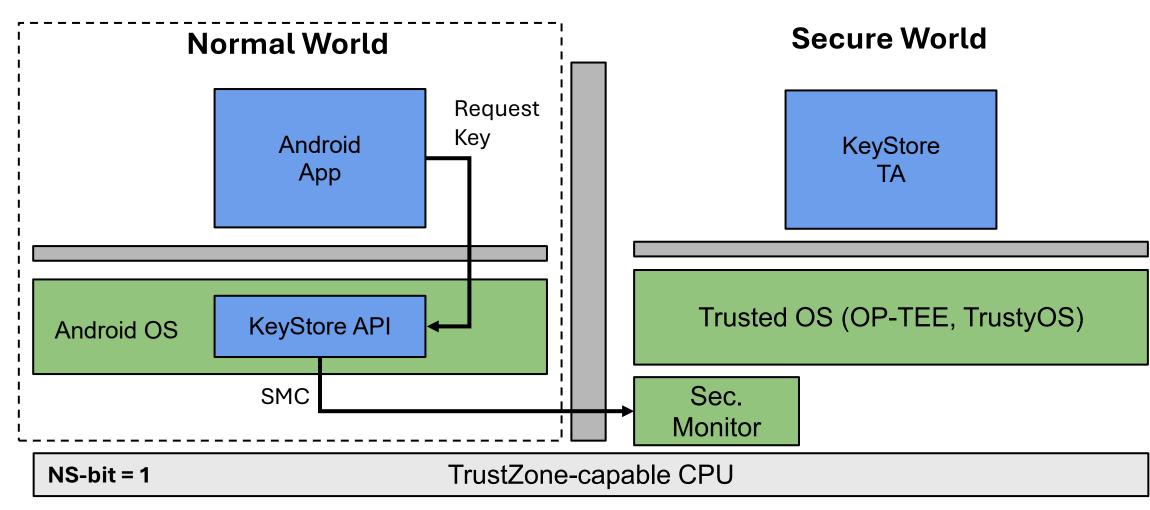
Android App will call the KeyStore API to request a key



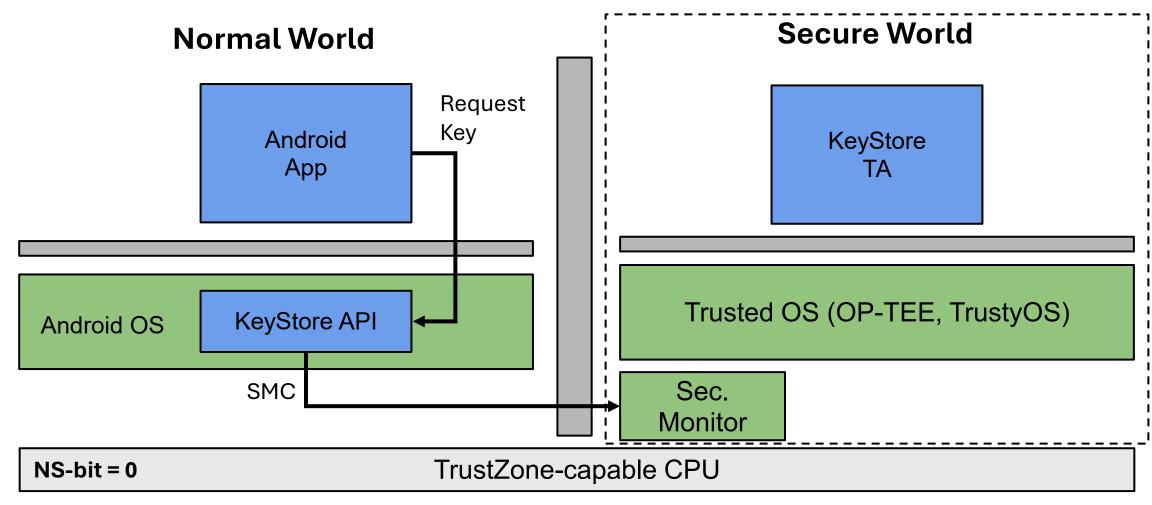
KeyStore API forwards the request to TEE How??



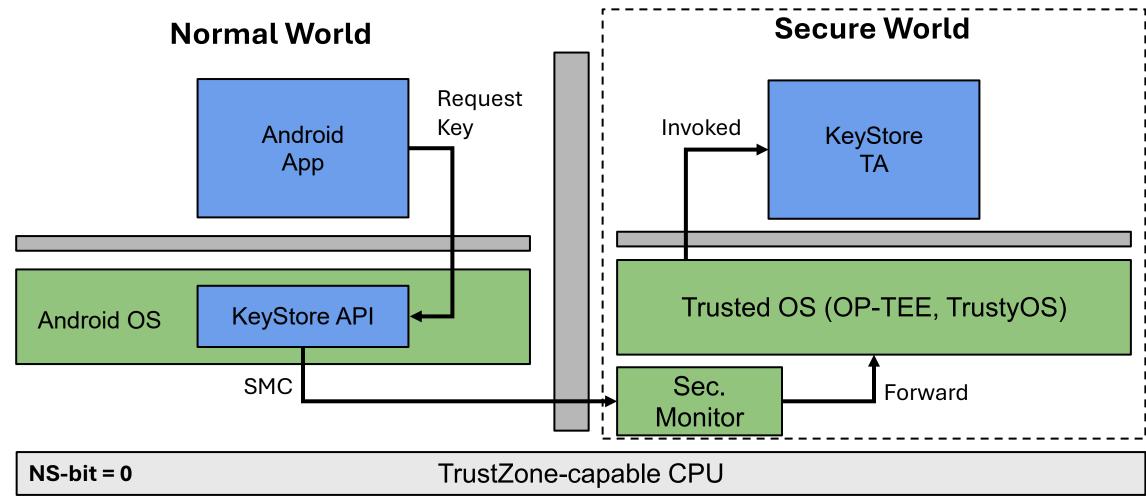
KeyStore API forwards the request to TEE How??



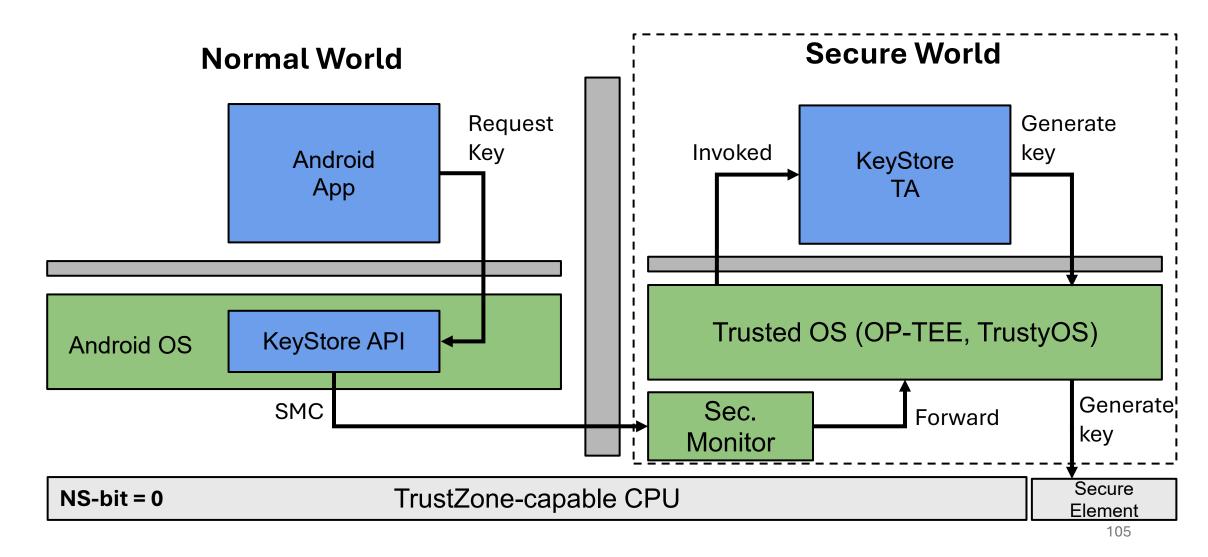
Security Monitor (atomically) switches from Normal World to Secure World



KeyStore TA is invoked. It then generates a key pair (e.g., AES, RSA) – how?



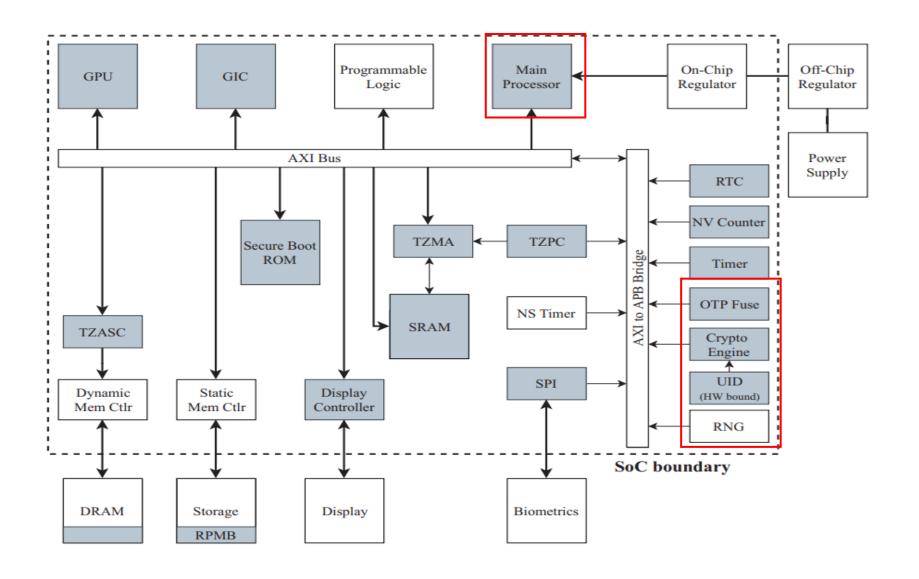
KeyStore TA is invoked. It then generates a key pair (e.g., AES, RSA) - how?



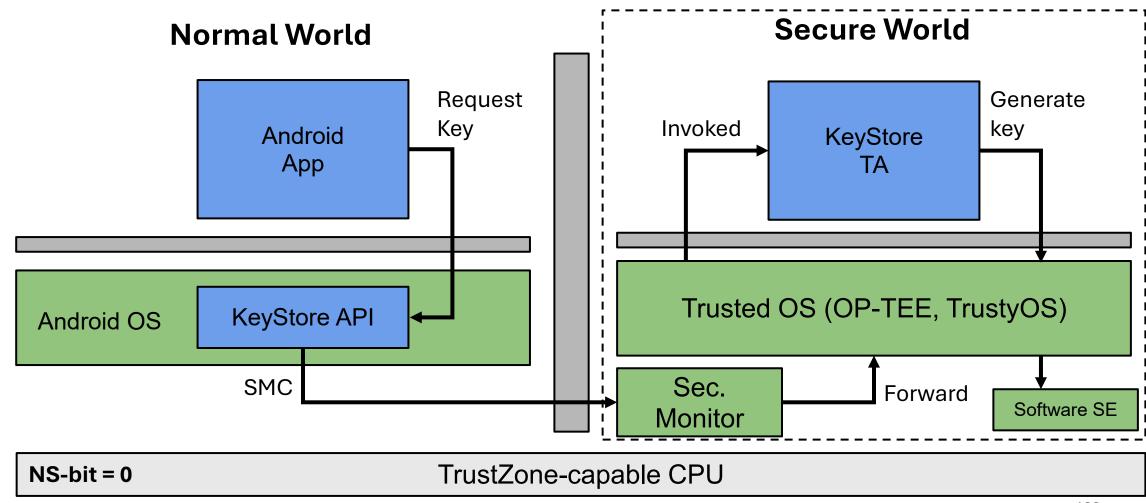
What is the Secure Element?

- Dedicated hardware module
 - Similar idea to TPM or HSM (sometimes is that)
 - Isolated component designed to handle cryptographic operations
- Very version-specific
 - Examples:
 - Separate PUF-based logic outside the CPU core but in the SoC (ANI1271):
 - StrongBox

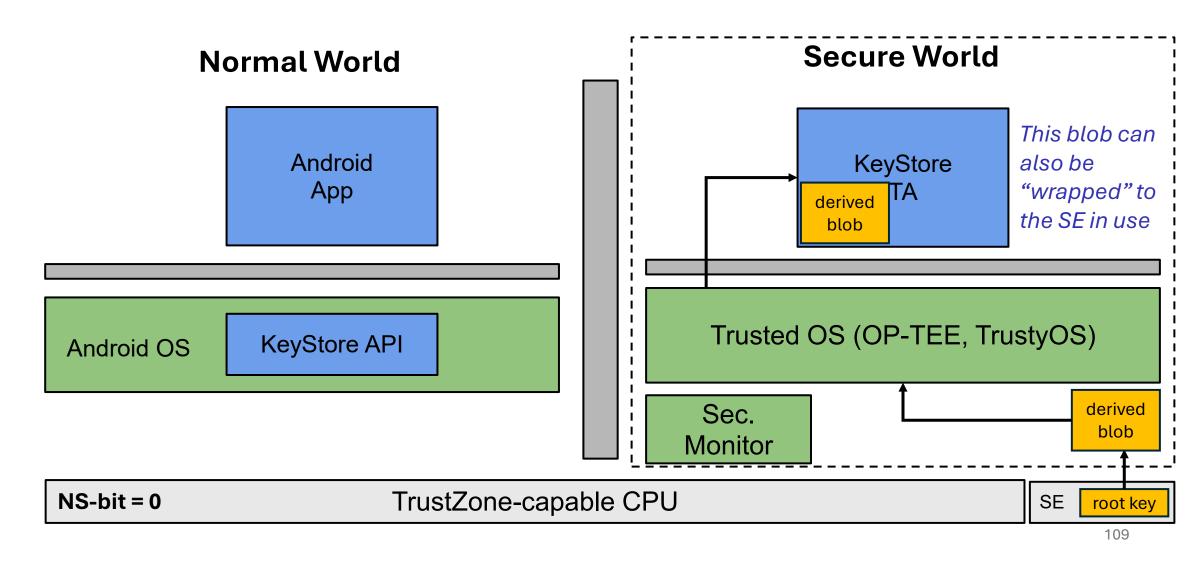
Recall:



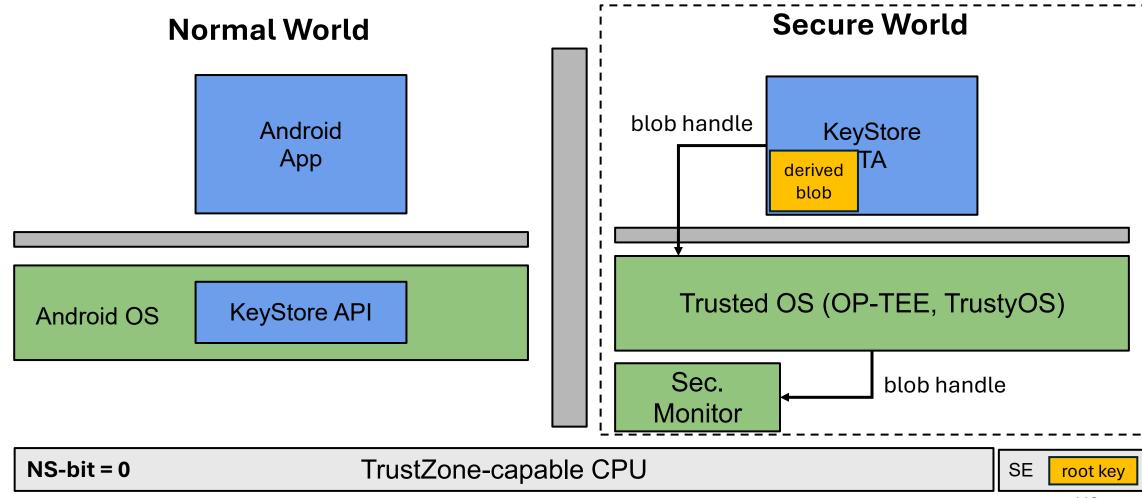
The TEE itself could be used for the SE, but less common



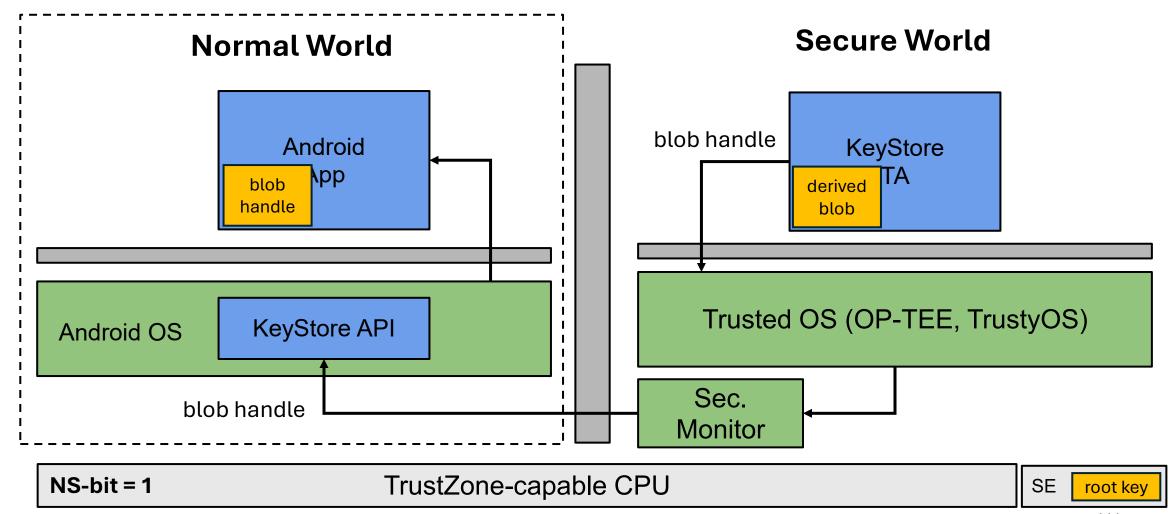
Generates key blob derived from a root key, stored in KS TA private memory



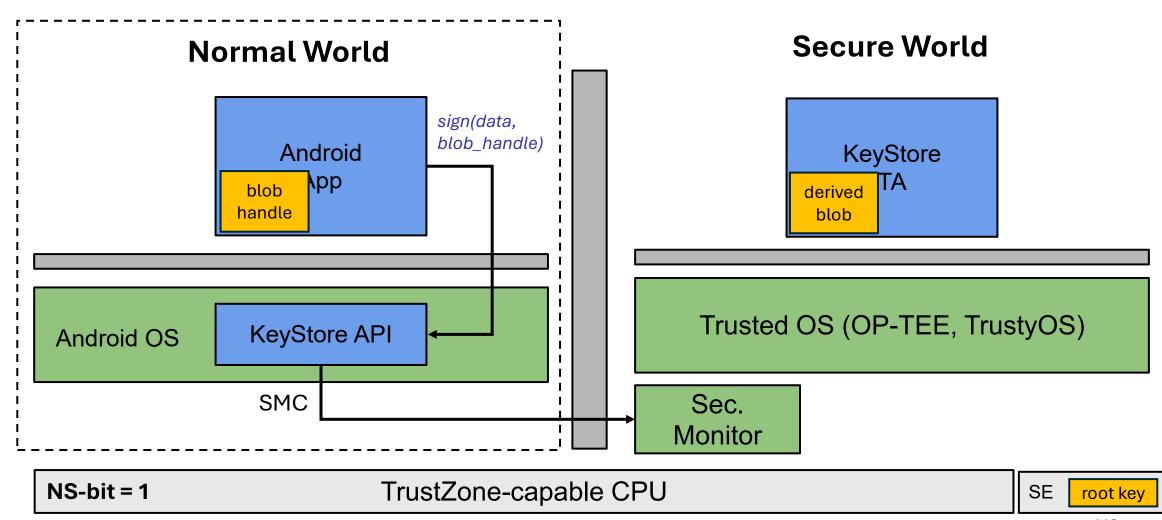
Returns a blob handle



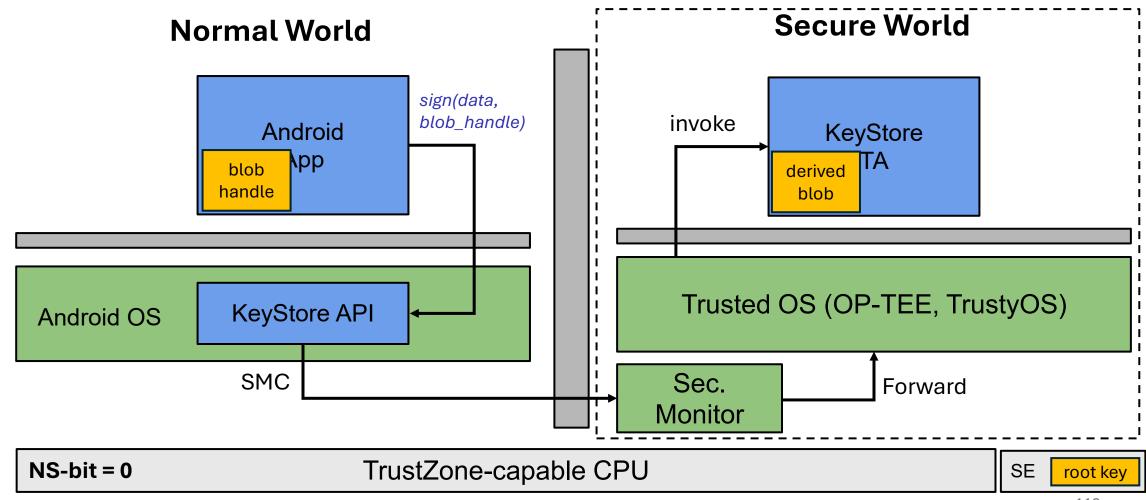
Atomic context switch, then store blob handle in Android App memory



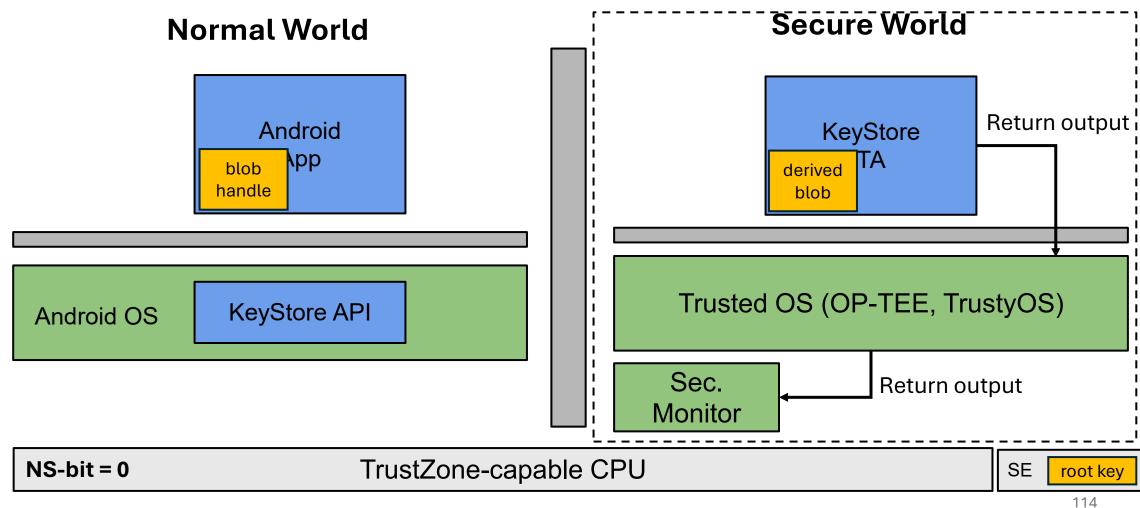
How about using the key? Invoke a KeyStore function, leading to SMC



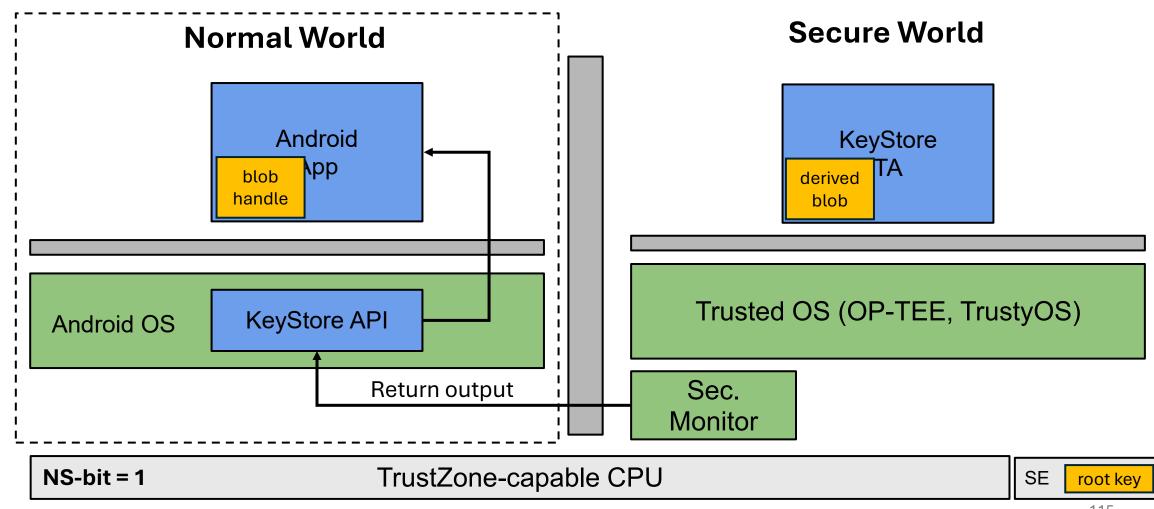
After being invoked, KeyStore TA retrieves key and performs operation



After being invoked, KeyStore TA retrieves key and performs operation



After being invoked, KeyStore TA retrieves key and performs operation



Other features of Android:

App signing

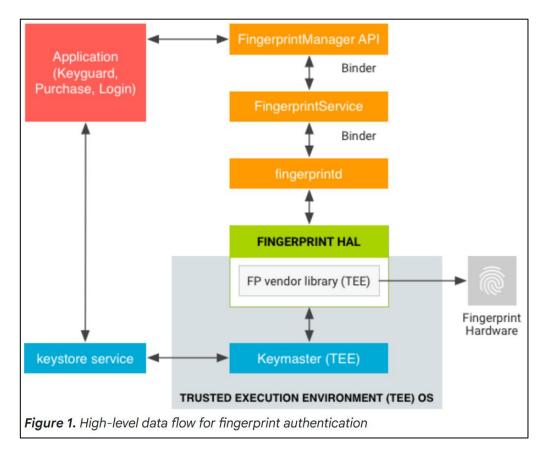
- Every app must be signed by the developers
- Unsigned apps are rejected by Google play or the package installer

Biometrics

- Part of tiered authentication model fingerprint senors
- Relies on the keystore for secure storage

Other features of Android:

Biometrics (continued)

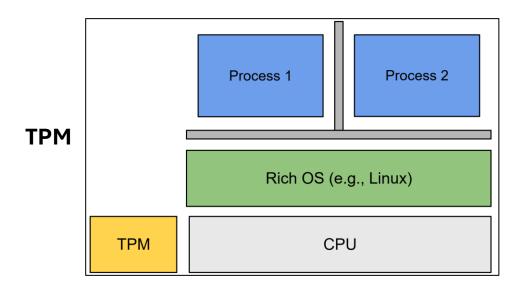


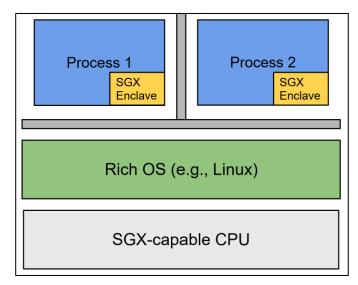
Other features of Android:

- App signing
 - Every app must be signed by the developers
 - Unsigned apps are rejected by Google play or the package installer
- Biometrics
 - Part of tiered authentication model fingerprint senors
 - Relies on the keystore for secure storage
- Verified Boot
- Rollback prevention
- Usable security "Private Space" sandboxed space with separate install of app

Closing thoughts

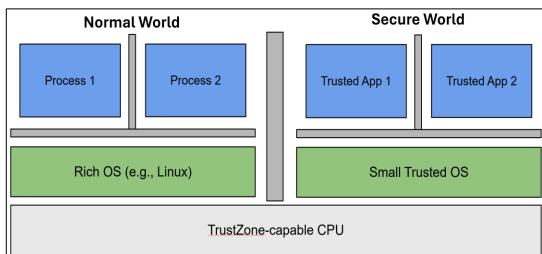
Various hardware security paradigms:





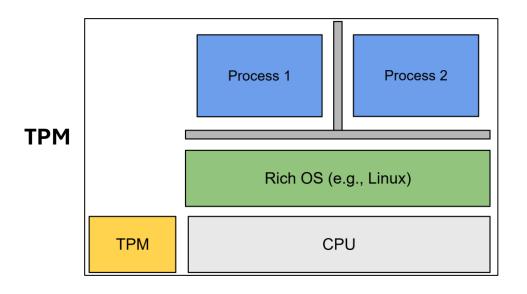
Intel SGX

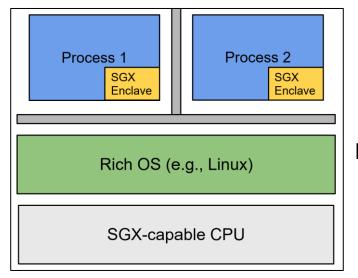




Closing thoughts

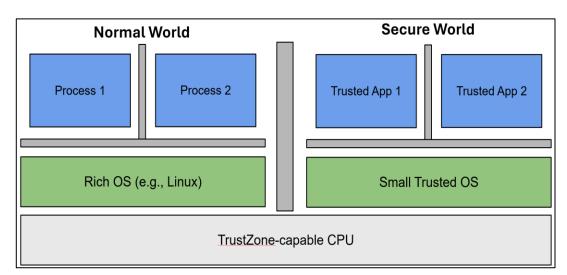
Various hardware security paradigms:





Intel SGX

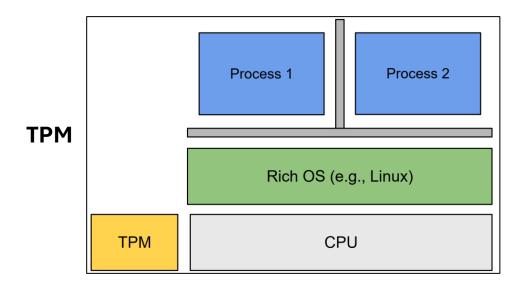


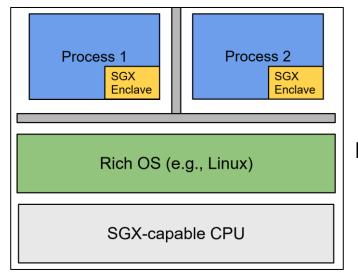


We saved the World!

Closing thoughts

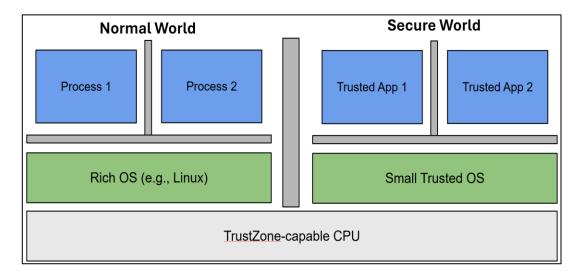
Various hardware security paradigms:





Intel SGX

ARM TrustZone



We saved the World!

Hopefully, there aren't any problems with these designs....

That's all for today!

Coming up....

Attacks on TPMs and TEEs

Reminders:

- A4 is due on July 25
- Research project proposal

That's all for today!

Resources:

- "Demystifying Arm TrustZone" great one!
- "TrustZone Explained: Architectural Features and Use Cases"
- ARM Docs on TrustZone-A
- Android security resources
- Android KeyStore
- HSE & SoC as SE
- "Safeguarding Crytographic Keys TEE and Strongbox in Android"
- "Mobile Platform Security"