CS 453/698: Software and Systems Security

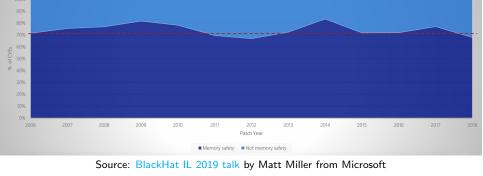
Module: An In-depth Study of Memory Errors Lecture: (casual discussion) memory-safe practices

> Meng Xu (University of Waterloo) Spring 2025

Outline

- Re-visit the statistics
- Memory-safe languages
- Capability Hardware Enhanced RISC Instructions (CHERI)

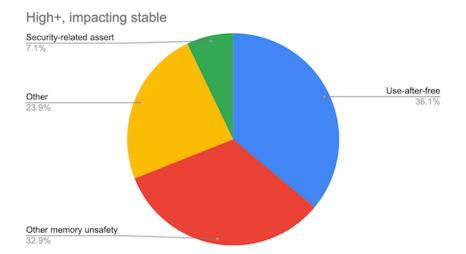
Memory errors are prevalent



% of memory safety vs. non-memory safety CVEs by patch year

Around 70% of all the vulnerabilities in Microsoft products addressed through a security update each year (2006 - 2018) are memory safety issues

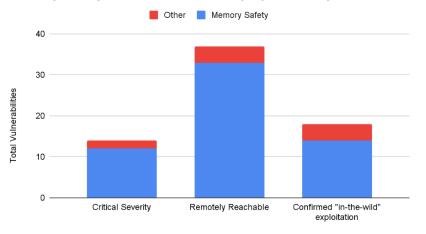
Memory errors are prevalent



Source: Chromium Memory Safety Report from Google.

Memory errors are prevalent

Memory Safety Vulnerabilities are Disproportionately Severe



Source: Blog post Memory Safe Languages in Android 13 from Google.

- Memory errors have universally accepted definitions (e.g., why the website is named Stack Overflow?)
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 - Once you find a memory error, you do not need to construct a working exploit to justify it
- Finding memory errors typically do not require program-specific domain knowledge (the bug is rooted in C/C++ language semantics instead of program logic)
 - If you have a technique that can find memory errors in one codebase, you can scale it up to millions of codebases developed in C/C++.

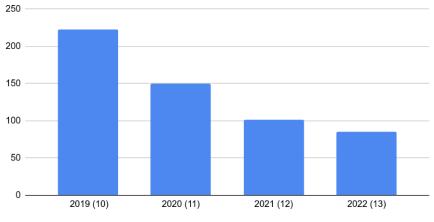
This is a personal note: one explanation why we have a disproportionately high number of memory errors reported amongst all security vulnerabilities is that — we know memory errors too well.

- Memory errors have universally accepted definitions (e.g., why the website is named Stack Overflow?)
 - Once you find a memory error, you do not need to diligently argue that this is a bug and not a feature
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 - Once you find a memory error, you do not need to construct a working exploit to justify it
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In fact, very few types of vulnerabilities meet these requirements.

Gradual adoption of memory-safe languages



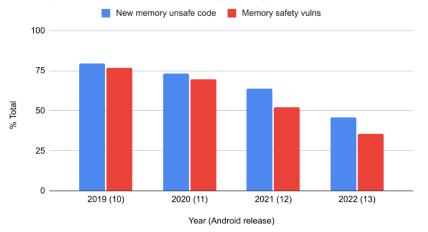


Year (Android release)

Source: Blog post Memory Safe Languages in Android 13 from Google.

Gradual adoption of memory-safe languages

Memory unsafe code and Memory safety vulnerabilities

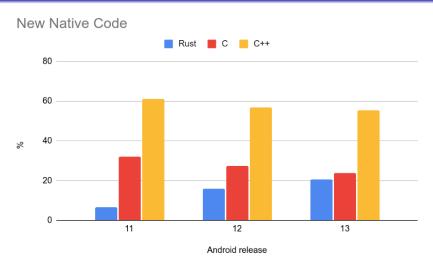


Source: Blog post Memory Safe Languages in Android 13 from Google.

Number of memory safety vulnerabilities correlates to the portion of unsafe code

Stats

Gradual adoption of memory-safe languages



Source: Blog post Memory Safe Languages in Android 13 from Google.

Rust on the rise in Android native implementations

Looking into the future

White House Press Release: Future Software Should Be Memory Safe on February 26, 2024.

ONCD Technical Report: Back to the Building Blocks: A Path Toward Secure and Measurable Software published in February 2024.



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- 2 Memory-safe languages
- 3 Capability Hardware Enhanced RISC Instructions (CHERI



(Potentially incomplete) list of memory-safe languages

Based on technical report "The case for memory-safe roadmaps" from NSA:

- C#
- Go
- Java
- Python
- Rust
- Swift

Java/Python

Q: How is spatial safety guaranteed?

Java/Python

Q: How is spatial safety guaranteed?

An example of the famous ArrayIndexOutOfBoundsException

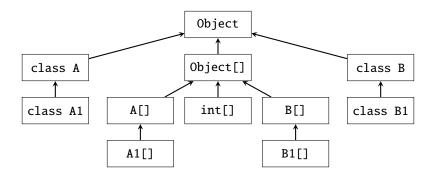
```
1 String[] names = { "tom", "bob", "harry" };
2 for (int i = 0; i \le names.length; i++) {
      System.out.println(names[i]):
4 }
```

The pattern looks similar to what happens in C when you have an out-of-bound memory access, but it is not a memory error in Java — Why?



How does Java VM track bounds

The key answer is: Java does not allow arbitrary casting.



- Upward cast is always allowed.
- Downward cast may be allowed.
- Re-interpret cast is never allowed.

Java/Python

Q: How is temporal safety guaranteed?



Java/Python

Q: How is temporal safety guaranteed?

- A: Garbage collection
- Automatically managed by the Java VM
- Identifies which objects are still in use (referenced) and which are not in use (unreferenced)
- Triggered upon certain conditions, such as
 - setting a reference to null
 - re-assigning a new object to a reference

Q: How is spatial safety guaranteed?

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A: Bounds check

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Q: How is temporal safety guaranteed?



Q: How is spatial safety guaranteed?

A: Bounds check

Q: How is temporal safety guaranteed?

A: Several ways, including

- Linear ownership transfer
- Lifetime annotation
- Reference counting

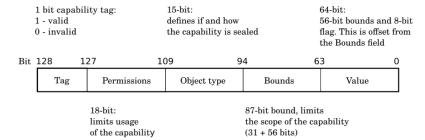
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Re-defining pointers

A pointer is not only an N-bit value representing a memory address, rather, it is a capability granting certain permissions to access a restrictive range in the memory address space.

CHERI memory capability



A "pointer", or rather, a memory capability, in the view of the CHERI Morello architecture (source of image: Pawel Zalewski's blog post).

```
#include <stdio.h>
int x=1;
int secret_key = 4091;
int main() {
   int *p = &x;
   p = p+1;
   int y = *p;
   printf("%d\n",y);
}
```

```
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Q: What will happen?

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```

```
x: signed int [@3, 0x14]

1

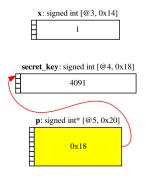
secret_key: signed int [@4, 0x18]

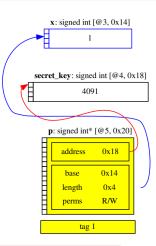
4091

p: signed int* [@5, 0x20]
```

Q: What will happen?

```
#include <stdio.h>
int x=1;
int secret_key = 4091;
int main() {
   int *p = &x;
   p = p+1;
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   printf("%d\n",y);
```





Q: What will happen?

CHERI software stack

Completely re-vamped software stack:

- Compilers: custom-made Clang/LLVM
- Operating systems: hand-tuned FreeBSD, FreeRTOS
- Applications: ported WebKit, OpenSSH, and PostgreSQL

 \langle End \rangle